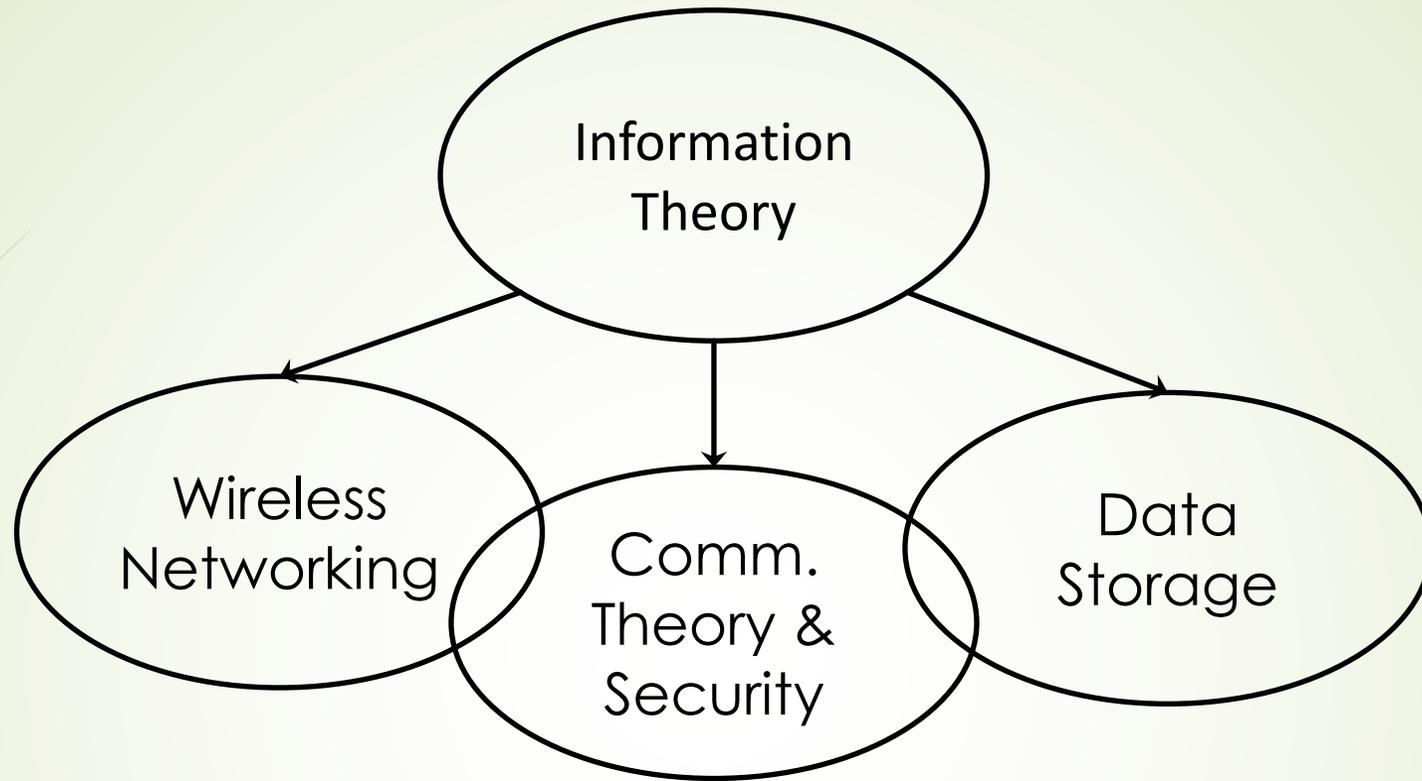


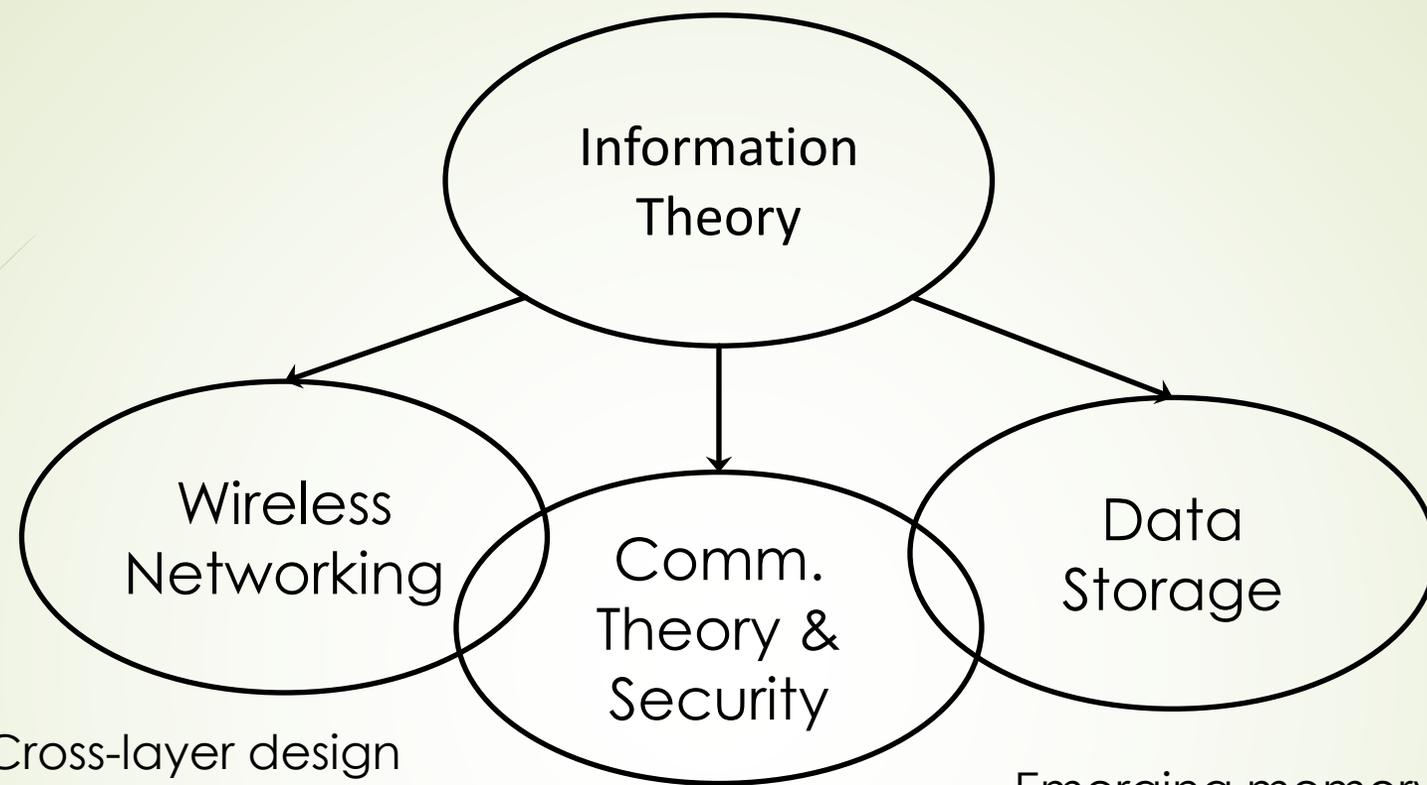


Defending WiFi Networks against Control Channel Attacks

Alireza Vahid, PhD

University of Colorado Denver

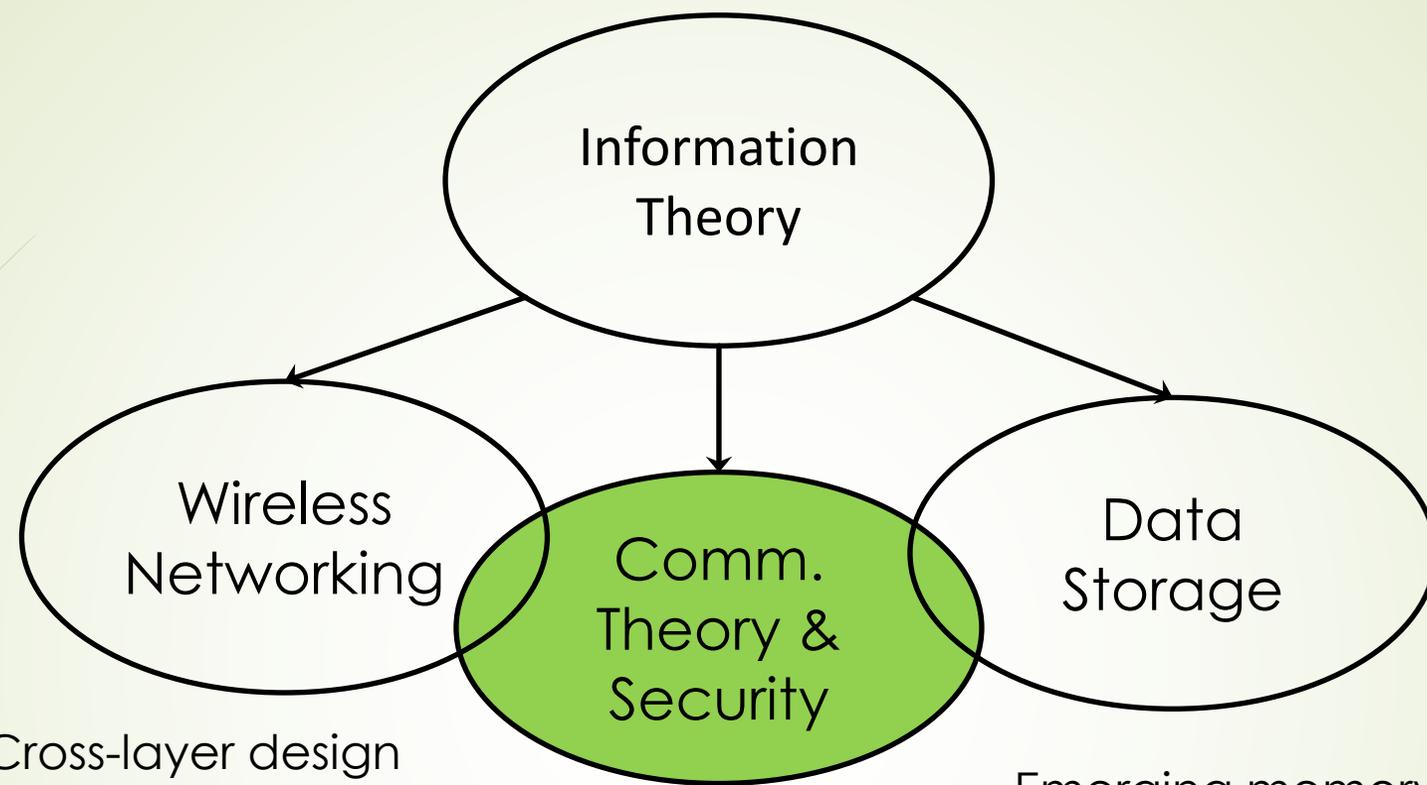




Cross-layer design
Spectrum sharing
Hardware-algorithm

Control ch. security
Feedback capacity

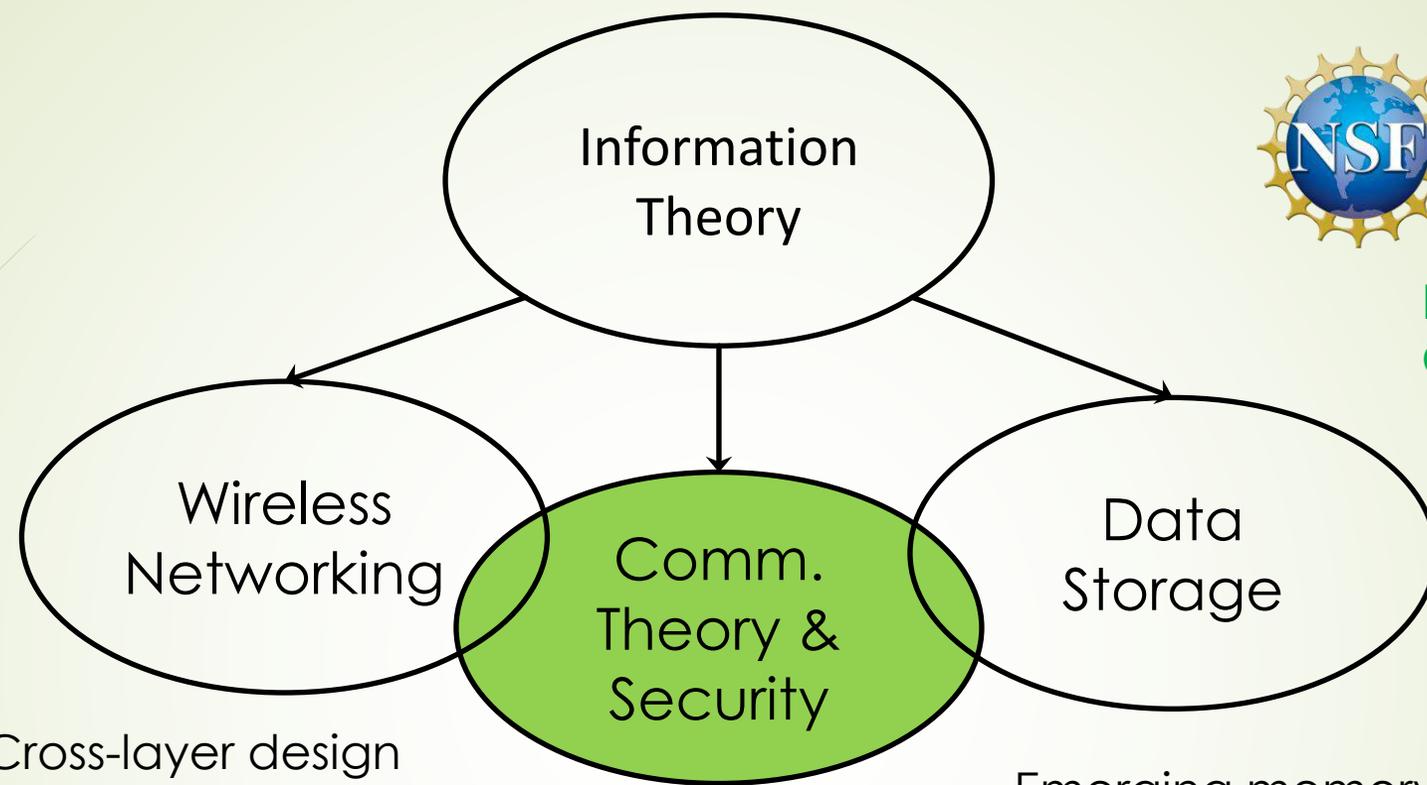
Emerging memory
Sequence assembly



Cross-layer design
Spectrum sharing
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Cross-layer design
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 Hardware-algorithm

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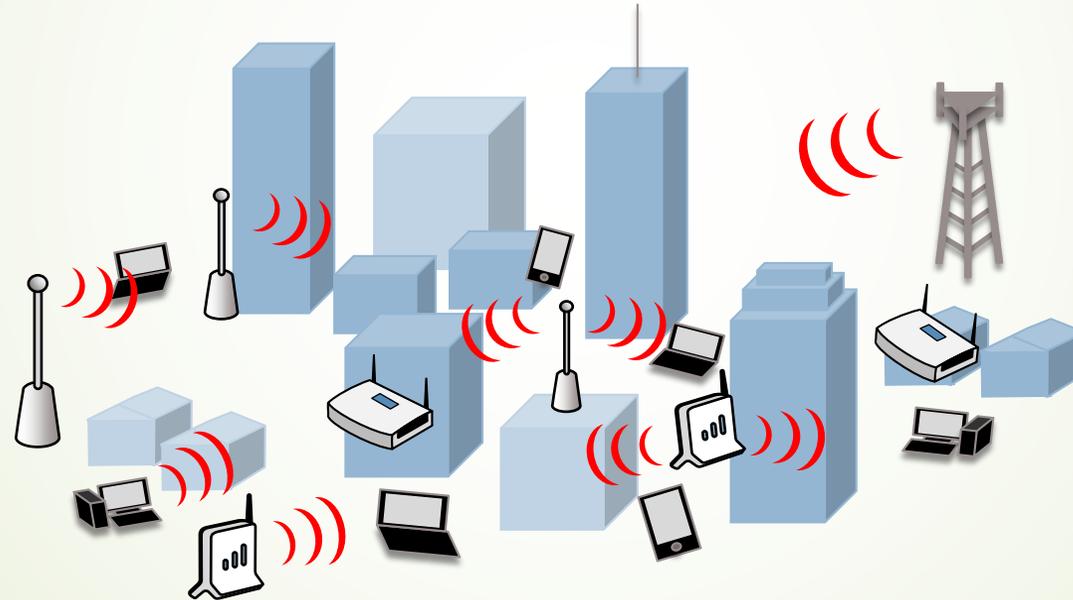


SWIFT
 NRDZ
 CNS

**Lab Venture
 Challenge Award!**



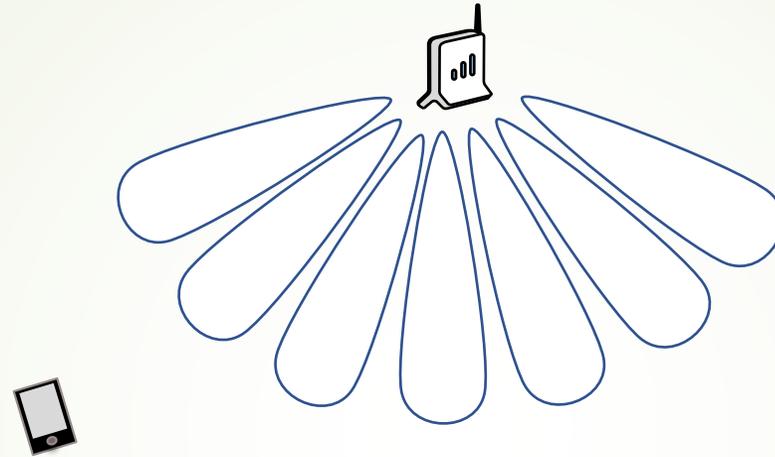
Future Wireless Networks



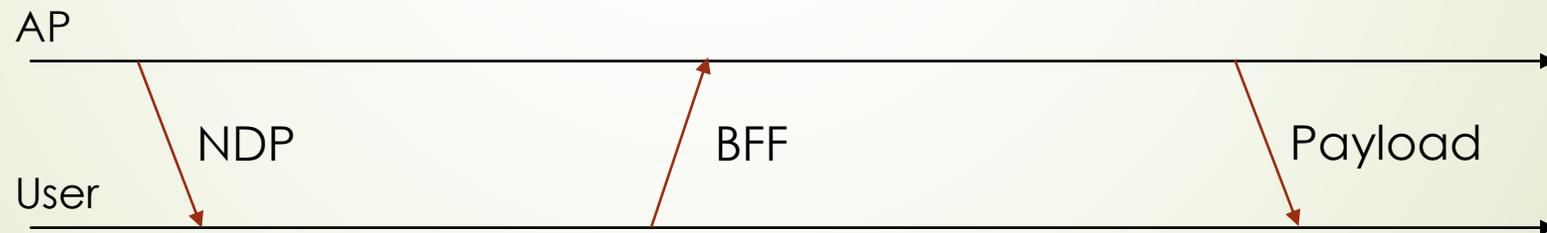
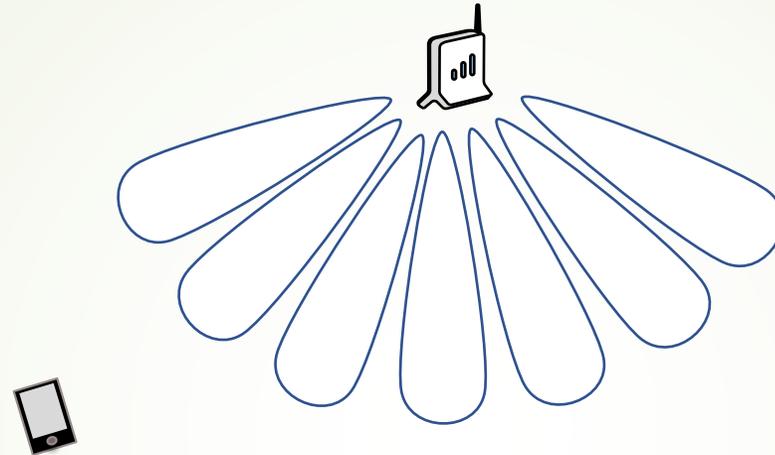
Future Wireless Networks



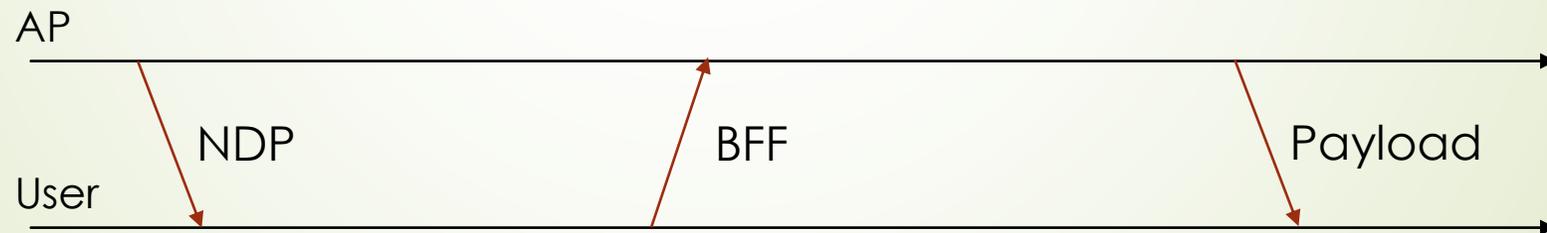
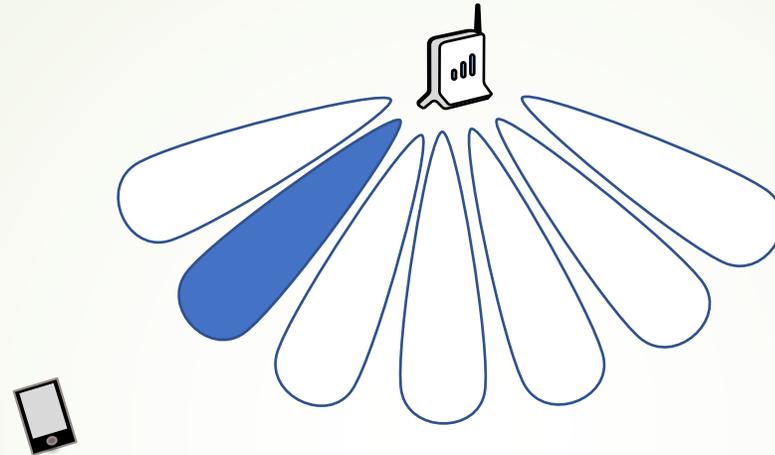
Control Channel Attack on WiFi-6 (-7)



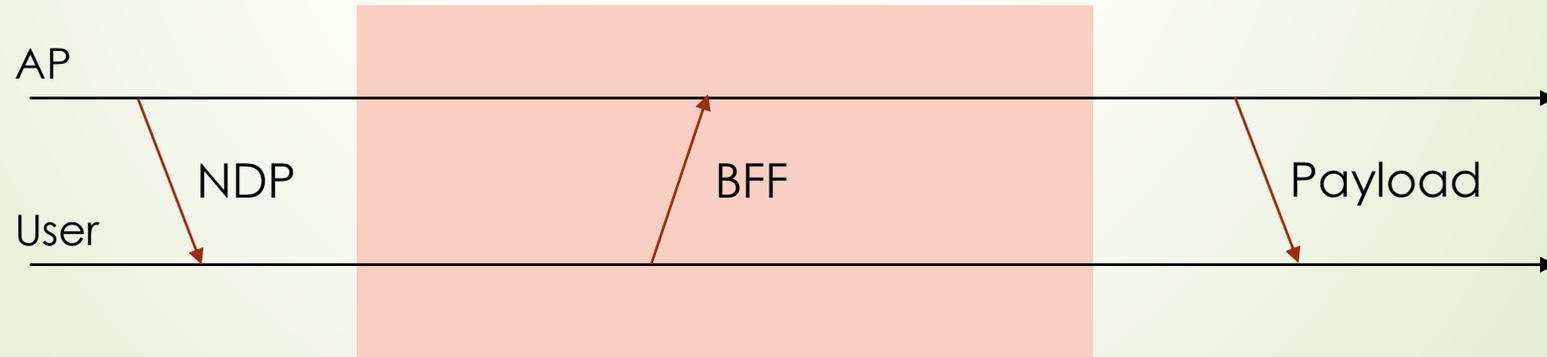
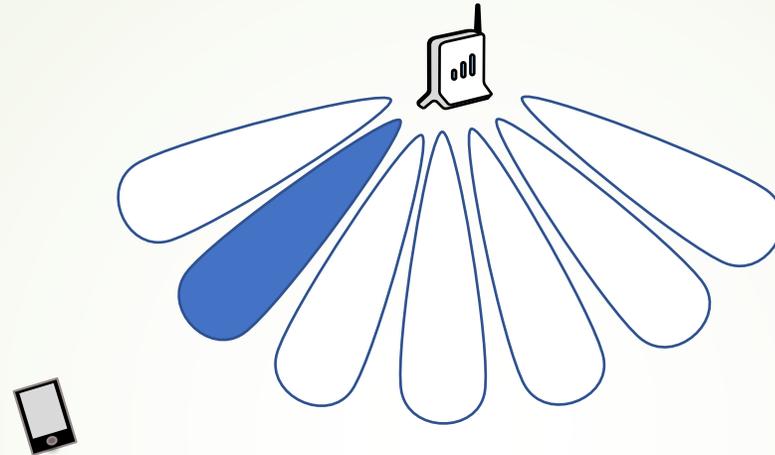
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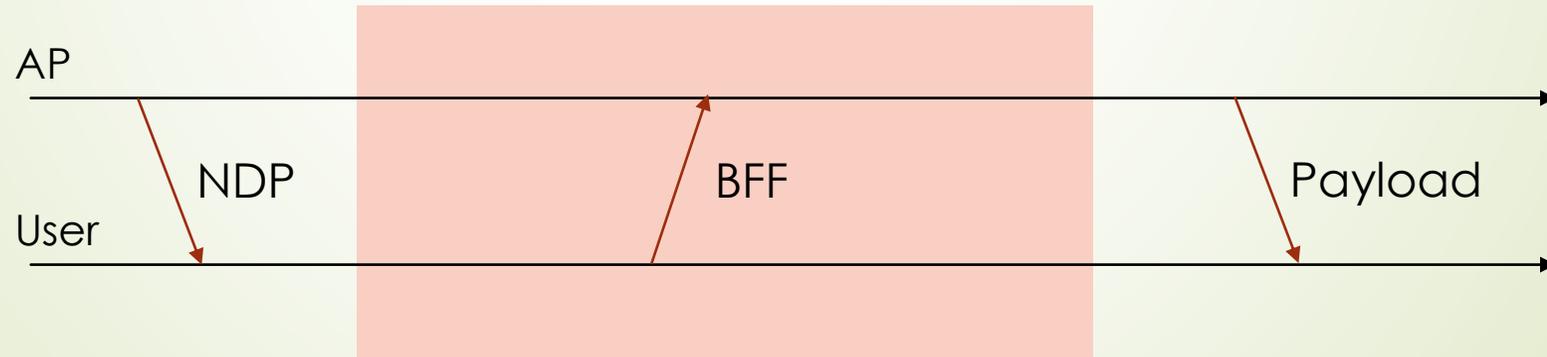
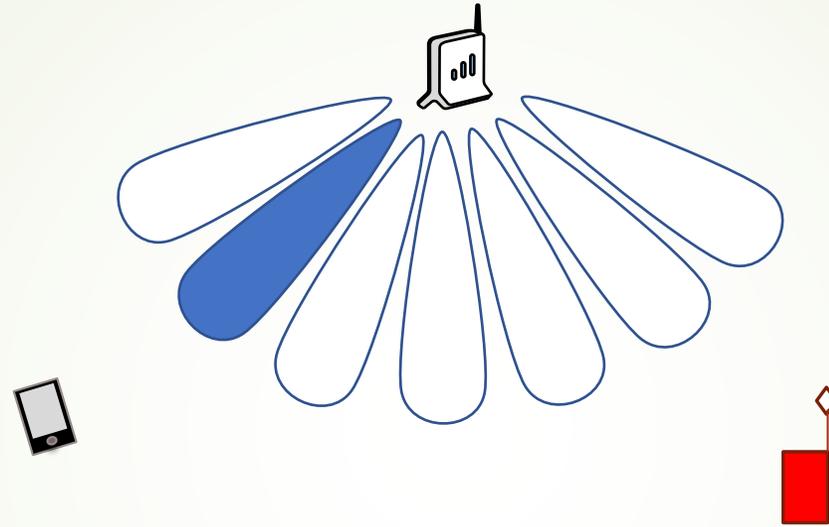
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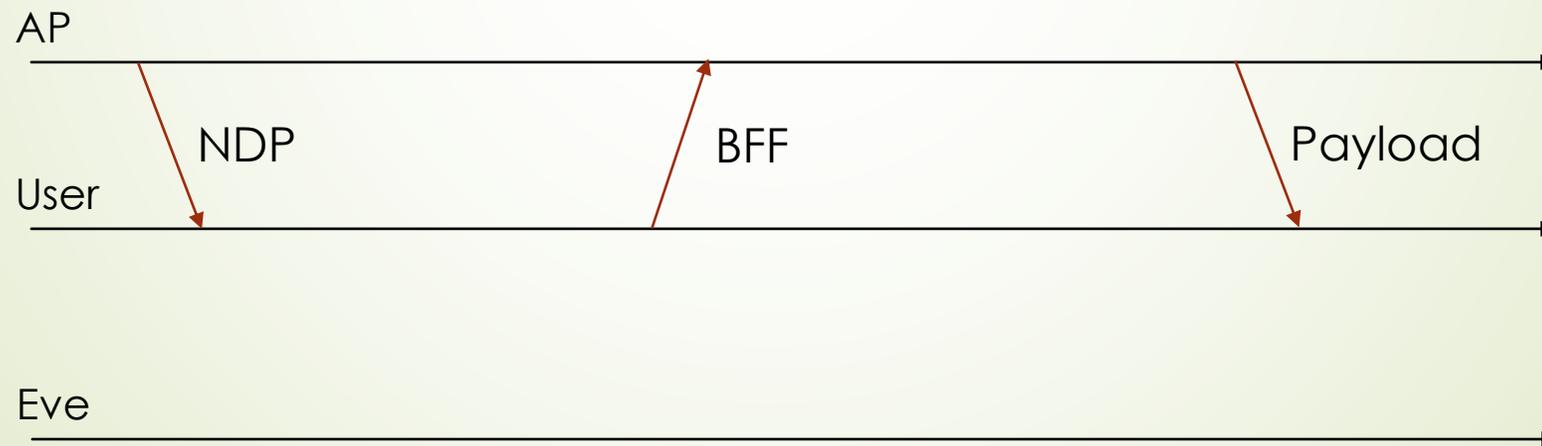
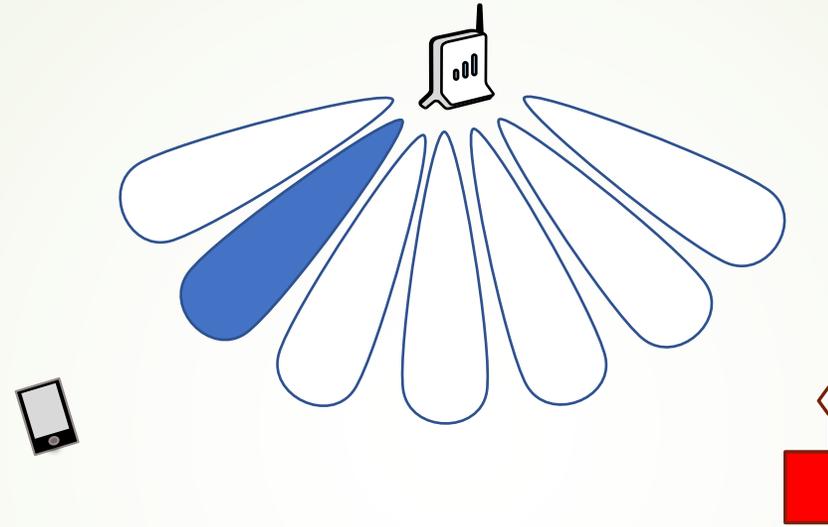
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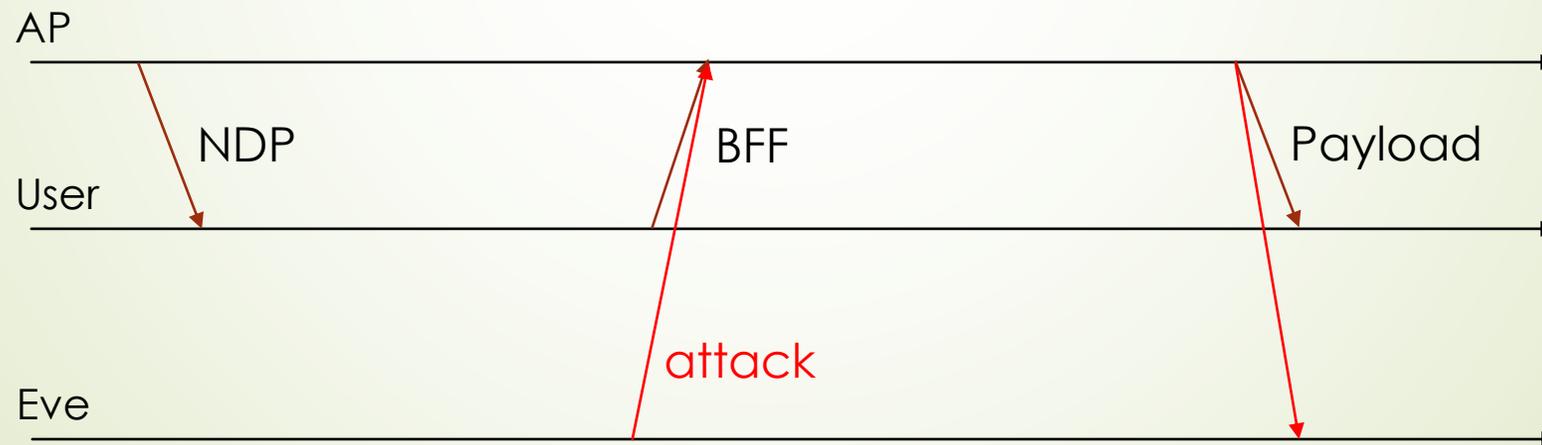
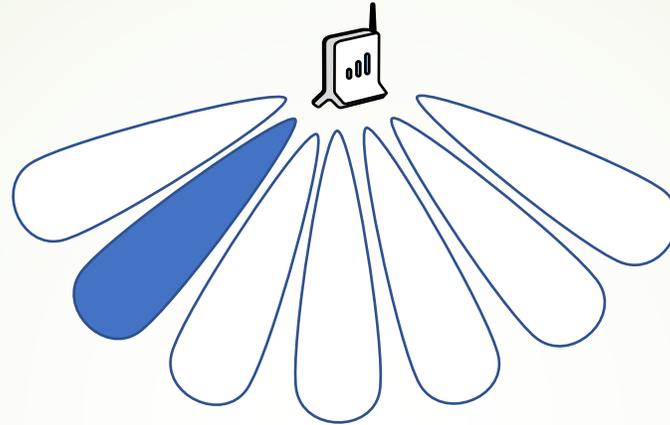
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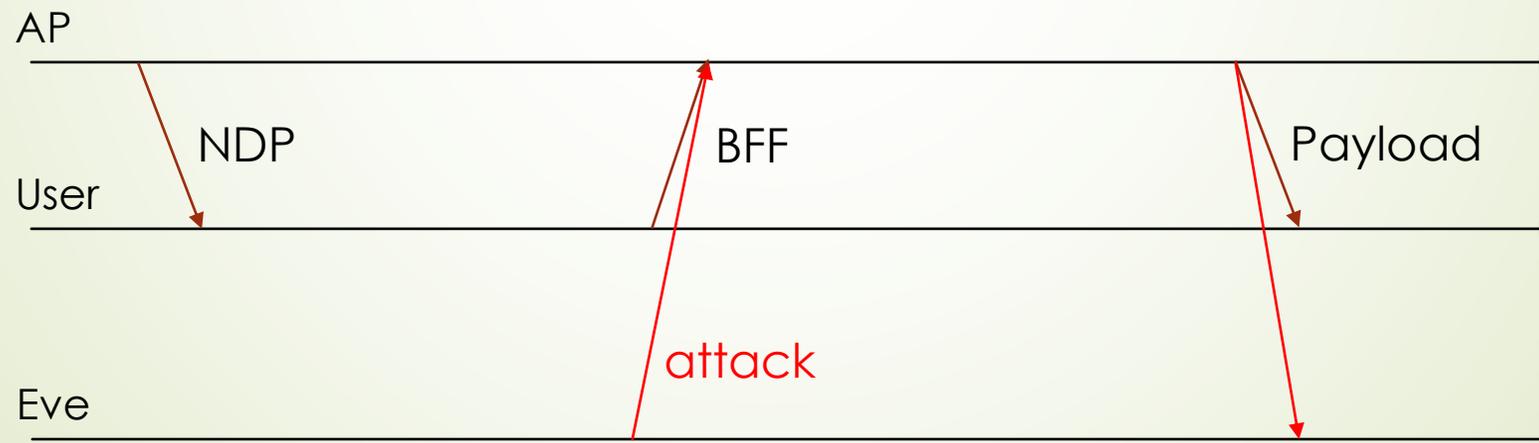
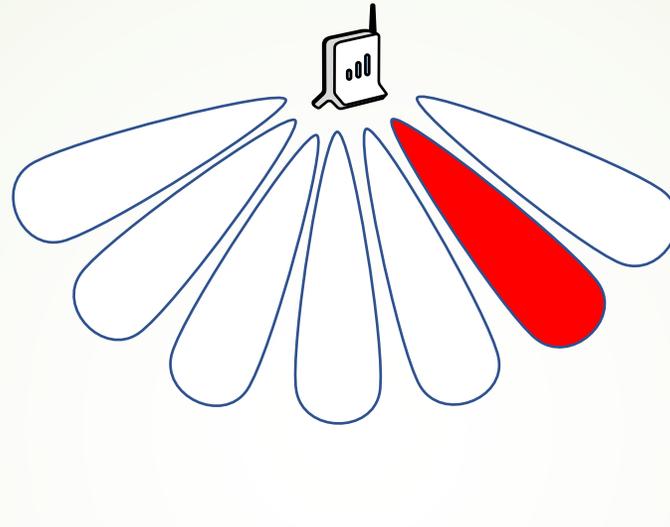
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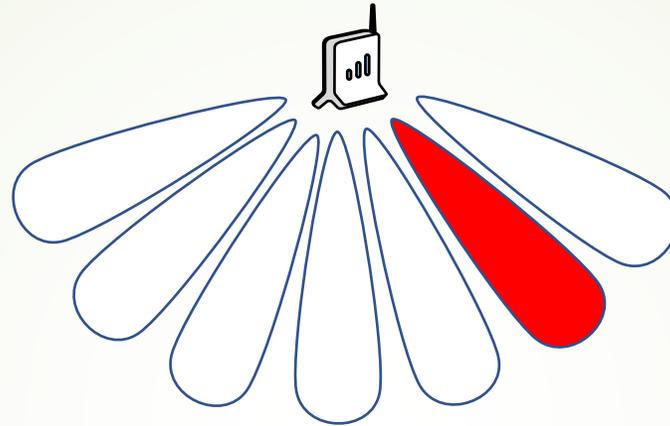
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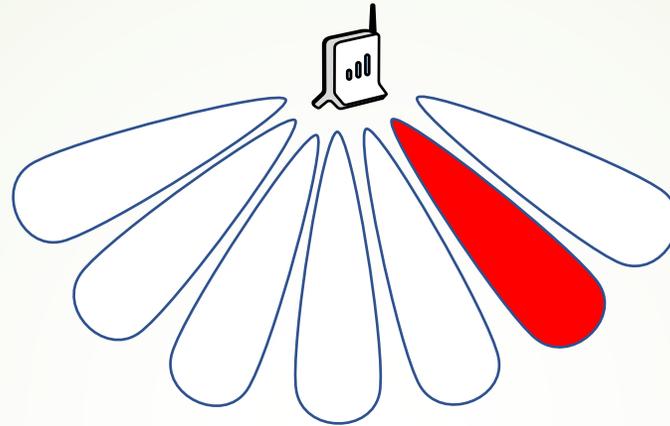
Control Channel Attack on WiFi-6 (-7)



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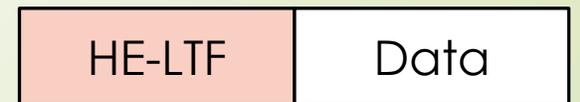
User



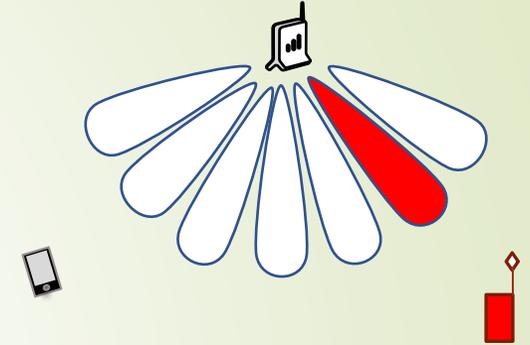
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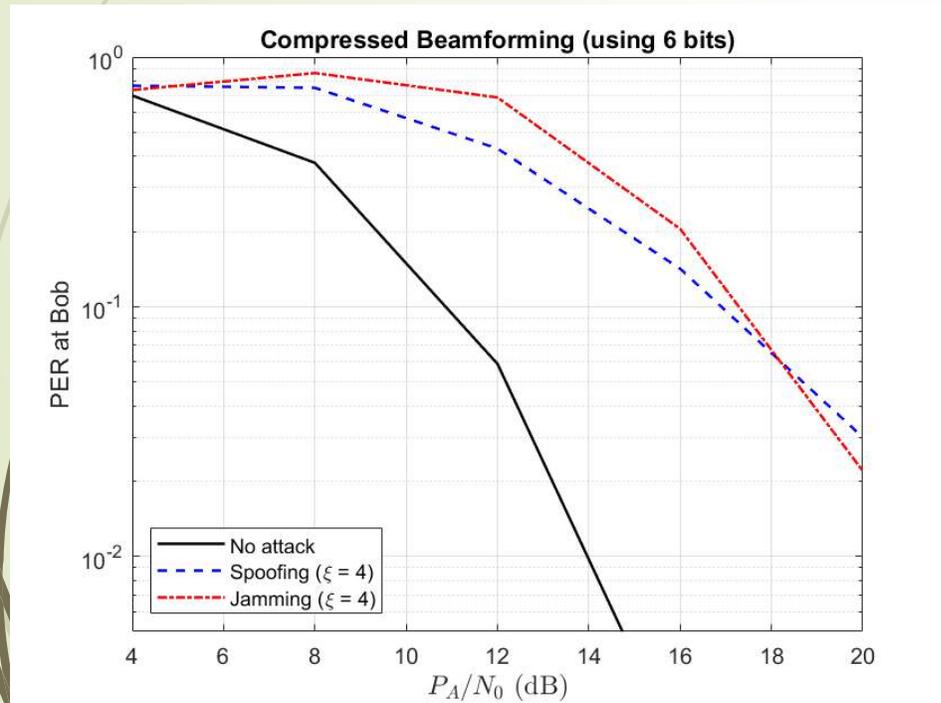
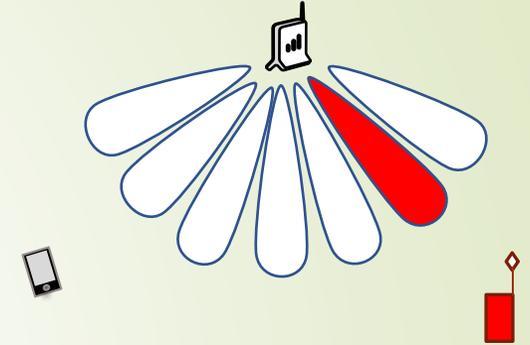
Eve



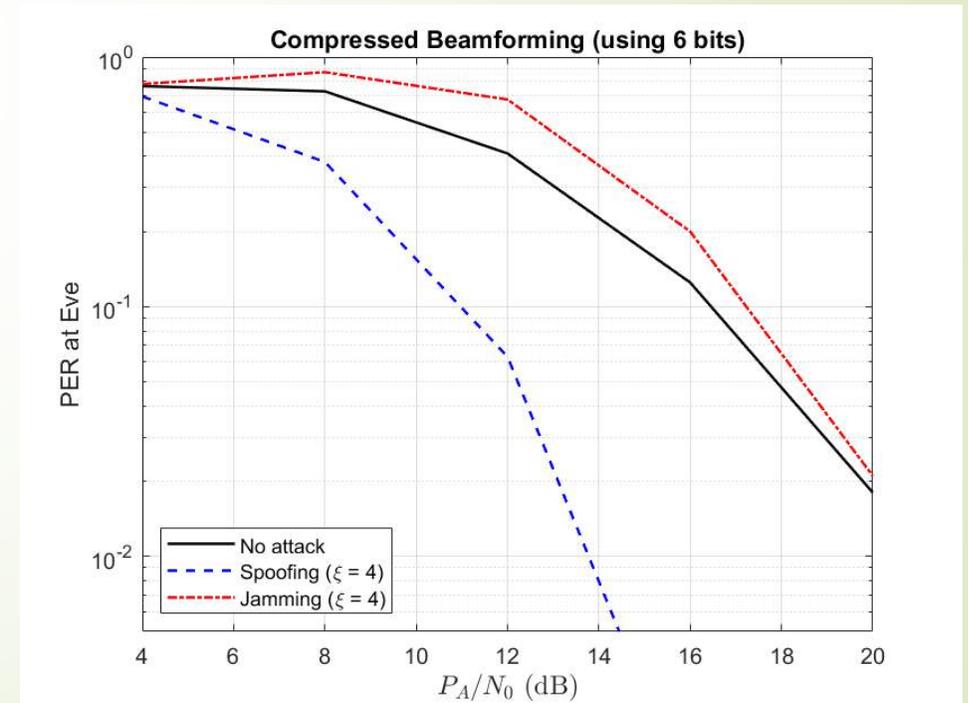
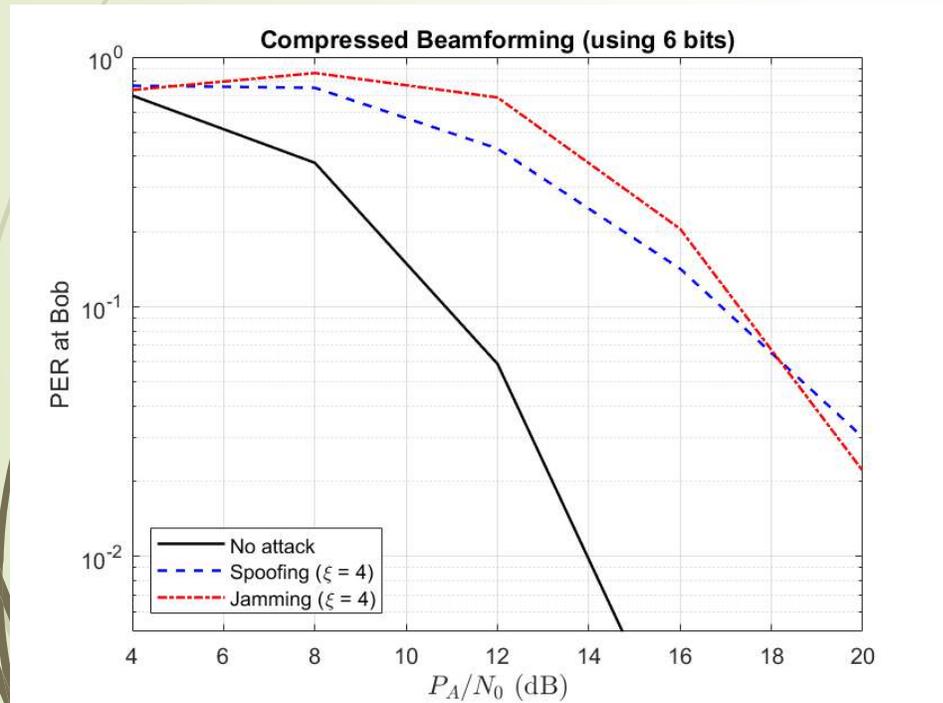
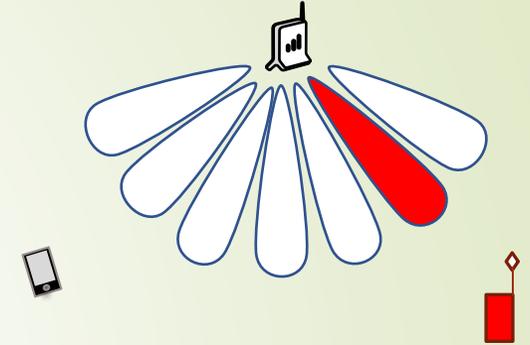
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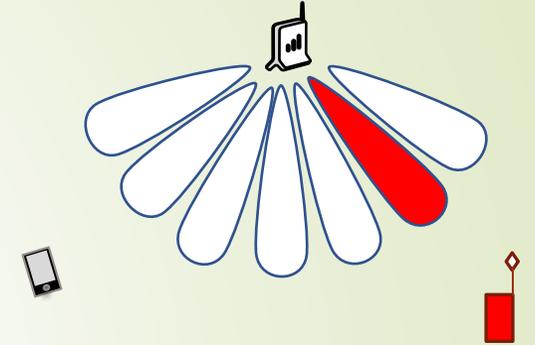
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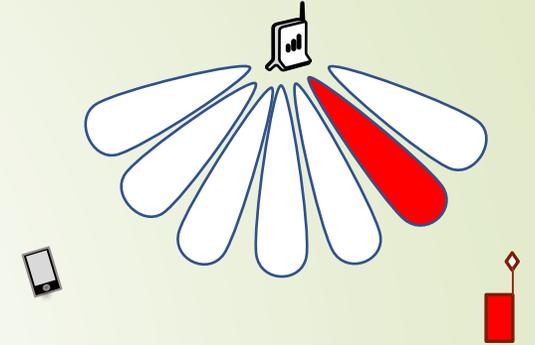
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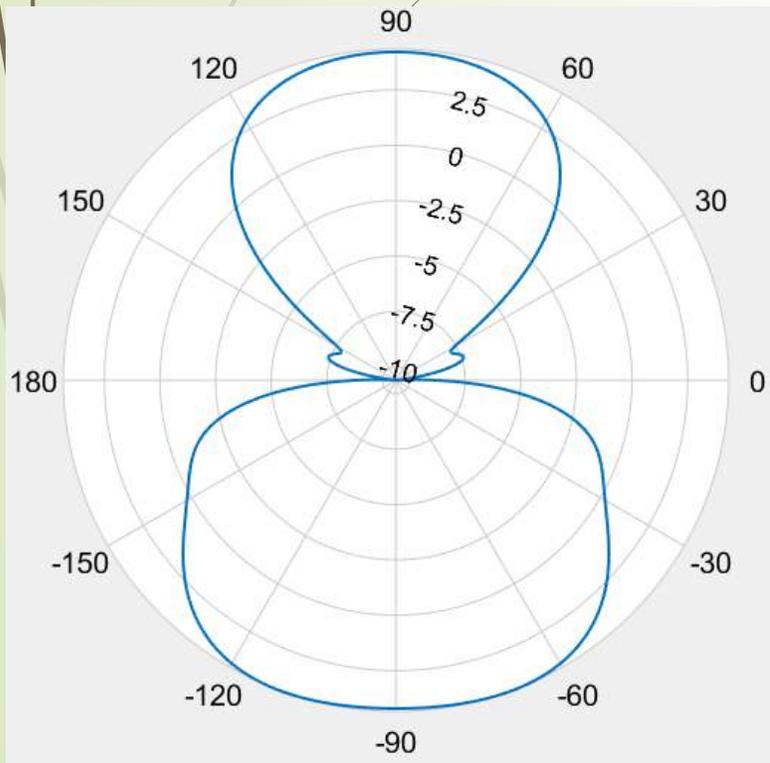
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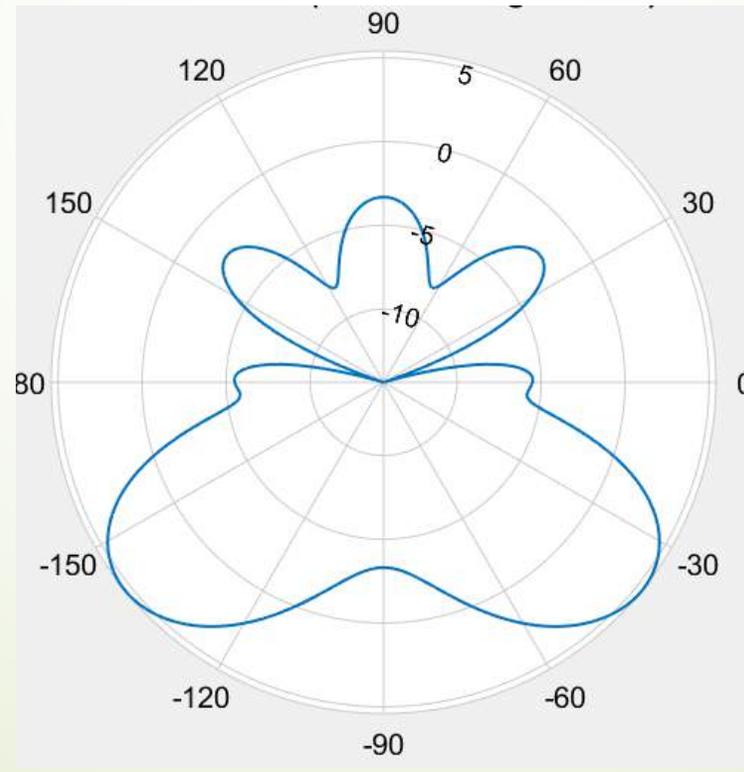
Control Channel Attack on WiFi-6 (-7)



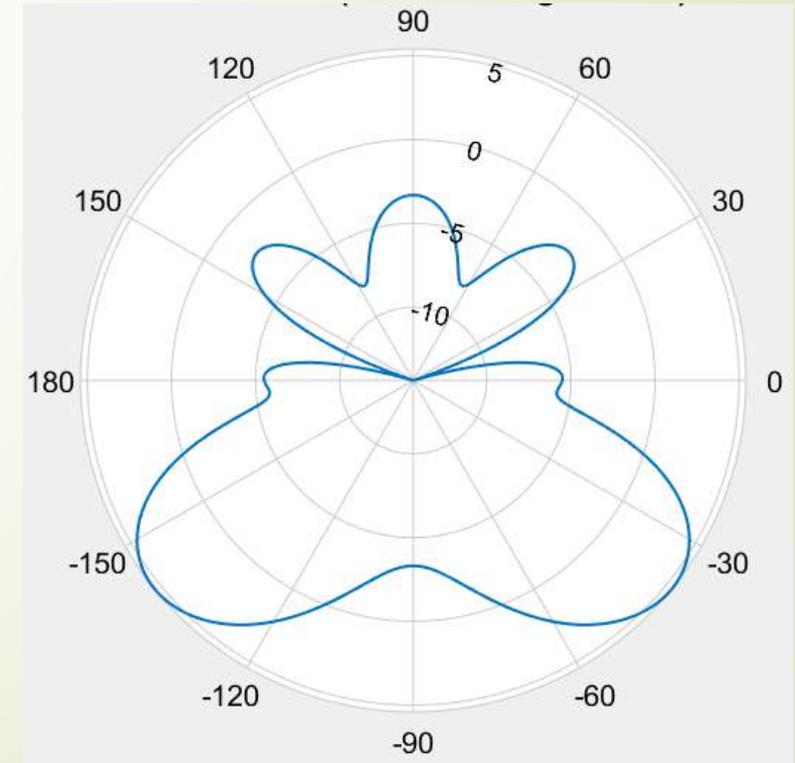
User



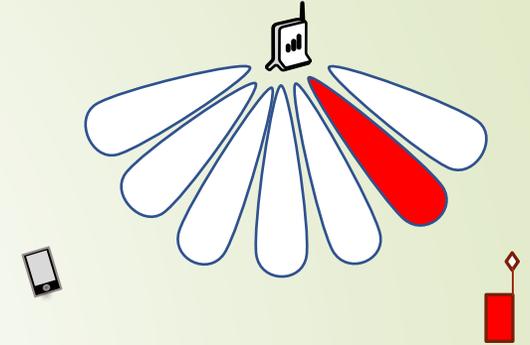
Eve



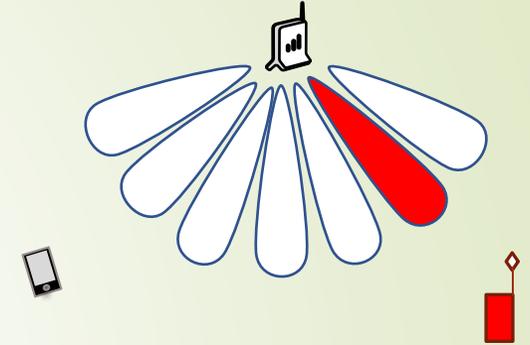
AP



Control Channel Attack on WiFi-6 (-7)

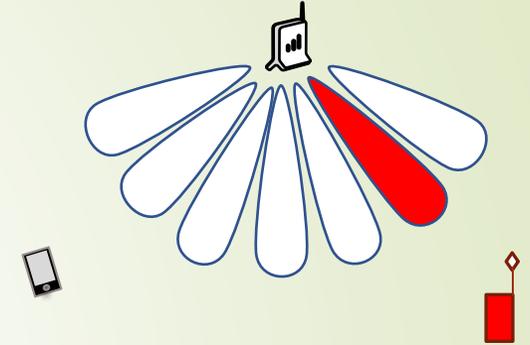


Control Channel Attack on WiFi-6 (-7)



- If the attack is successful, the following control packets (e.g., ACK/NACK) will be likely corrupted.

Control Channel Attack on WiFi-6 (-7)



- If the attack is successful, the following control packets (e.g., ACK/NACK) will be likely corrupted.
- Plan:
 - Quantify the impact
 - Detect
 - Defend

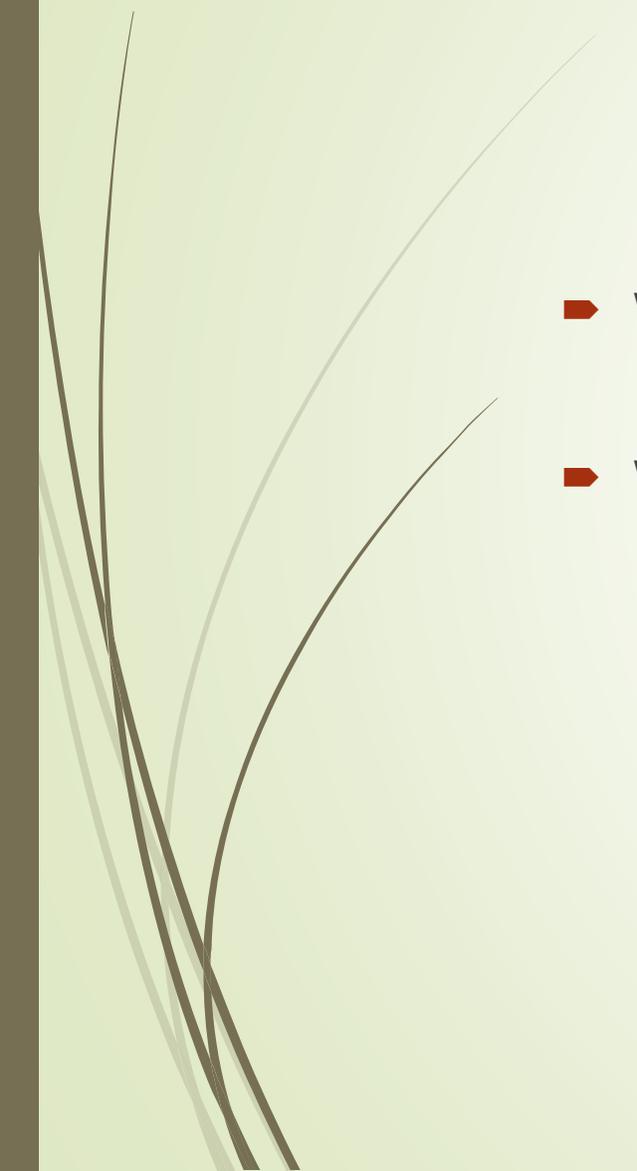


Understanding the fundamentals

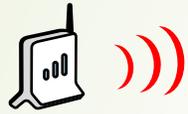
- ▶ We want to quantify the impact and devise the protocols accordingly.
- 



Understanding the fundamentals

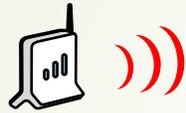
- ▶ We want to quantify the impact and devise the protocols accordingly.
 - ▶ We focus on small ACK/NACK control packets in a broadcast setting.
- 

Denial-of-service attack



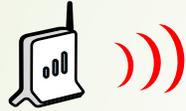
- We look at a fundamental model, the packet broadcast channel.

Denial-of-service attack



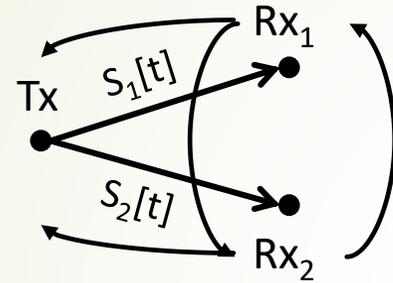
- ▶ We look at a fundamental model, the packet broadcast channel.
- ▶ Each user informs the transmitter whether the transmitted packet was received successfully or not.

Denial-of-service attack

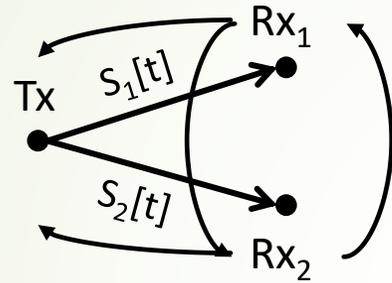


- ▶ We look at a fundamental model, the packet broadcast channel.
- ▶ Each user informs the transmitter whether the transmitted packet was received successfully or not.
- ▶ Why this model?

Channel model and baseline



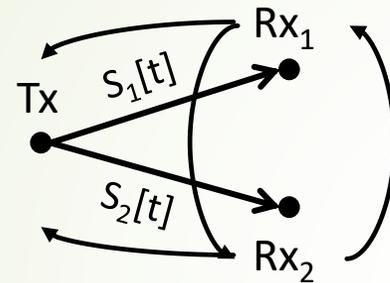
Channel model and baseline



$S_i[t]$ is Bernoulli ($1-\delta_i$)

$S_1[t]$ & $S_2[t]$ distributed
independently over time

Channel model and baseline

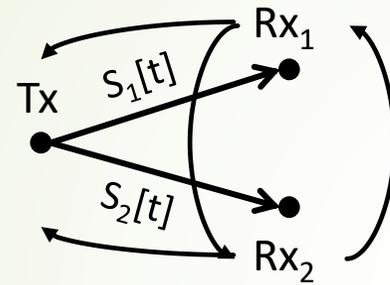


$S_i[t]$ is Bernoulli $(1-\delta_i)$

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Each Rx simply broadcasts its control packet

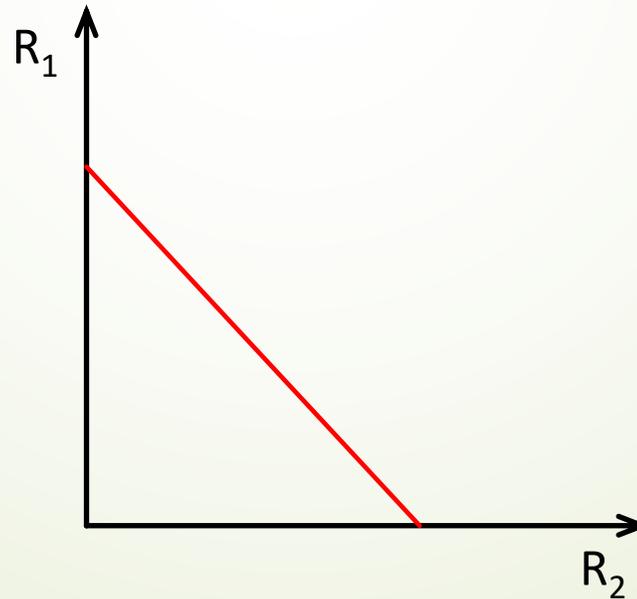
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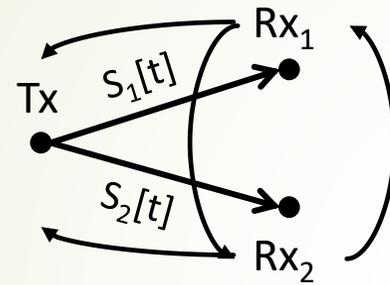
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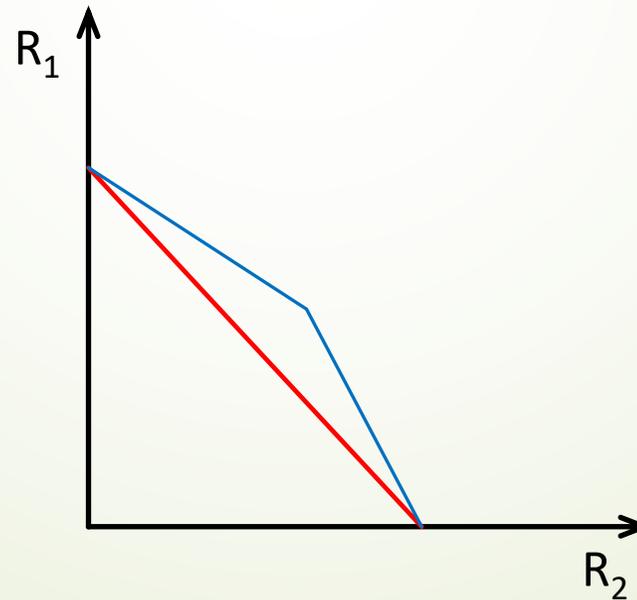
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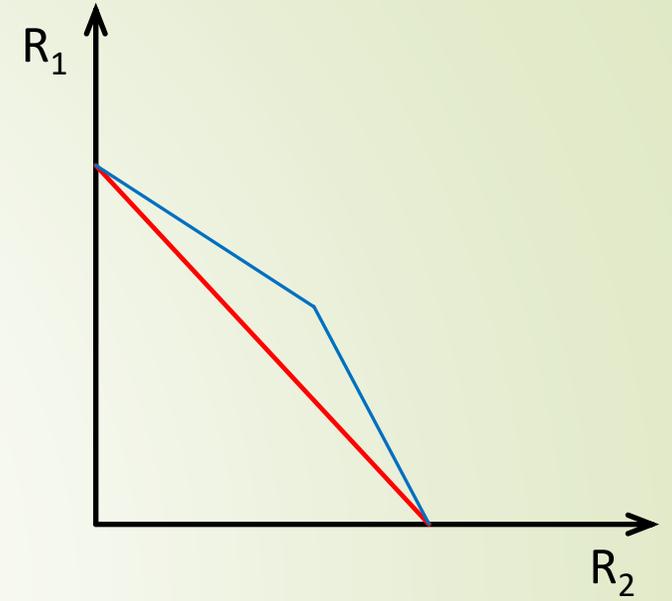
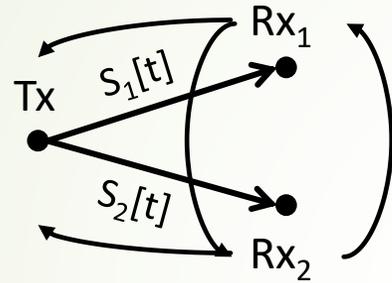
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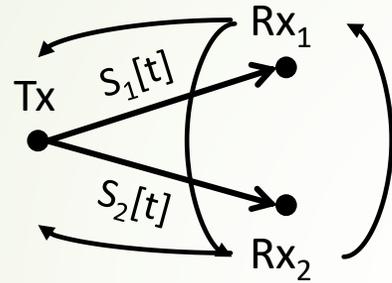
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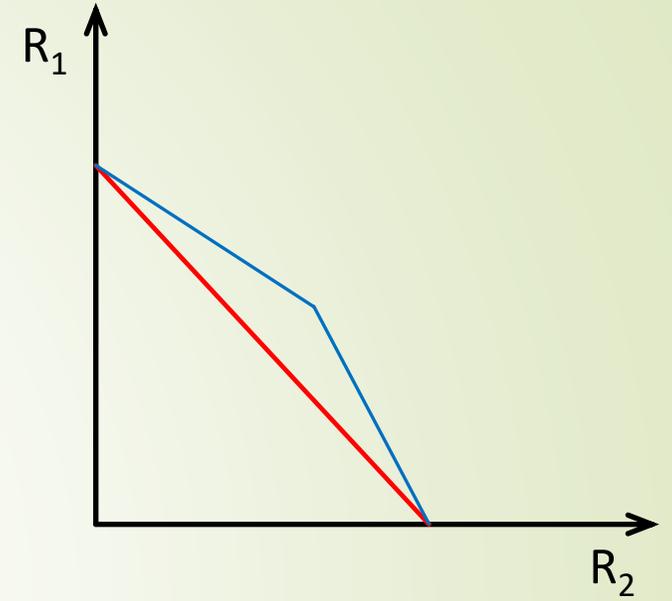
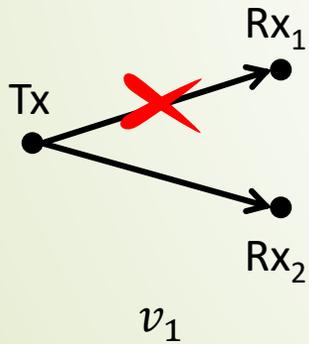
Protocol with no attack



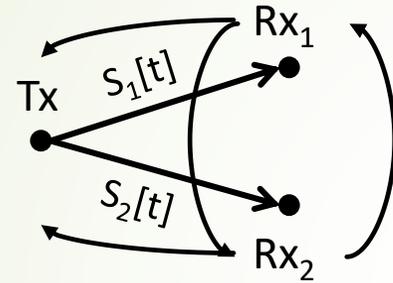
Protocol with no attack



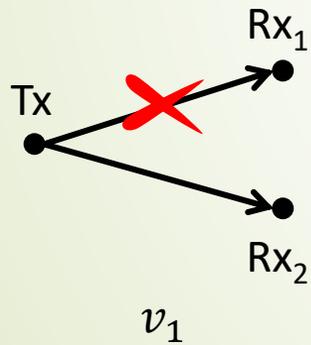
Send user 1's packets



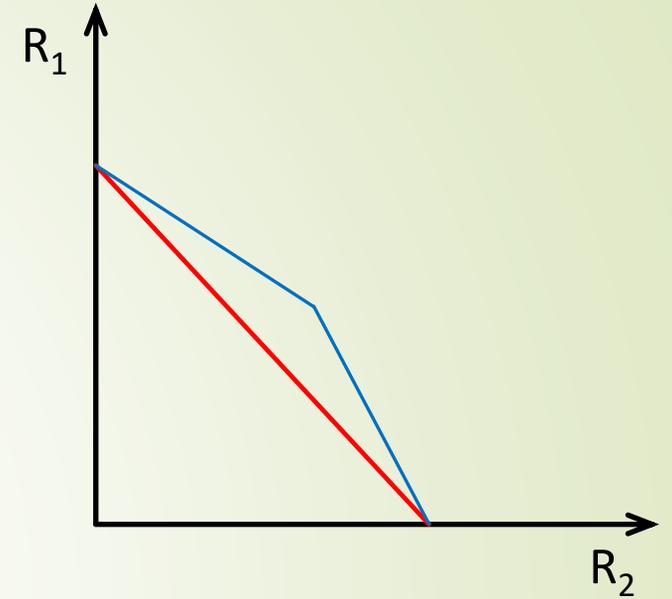
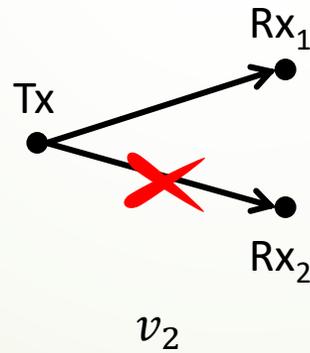
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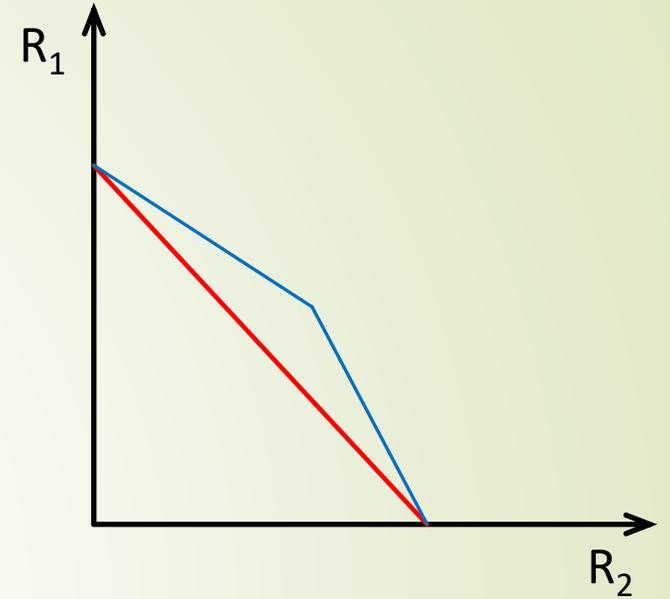
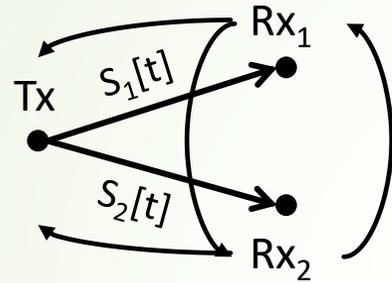
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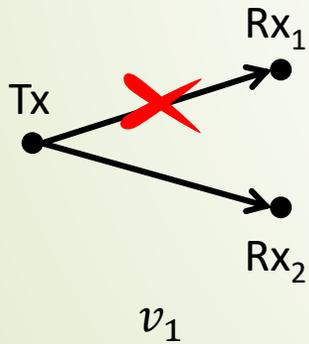
Send user 2's packets



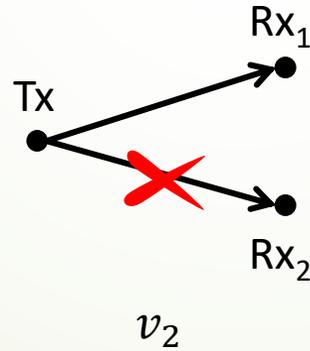
Protocol with no attack



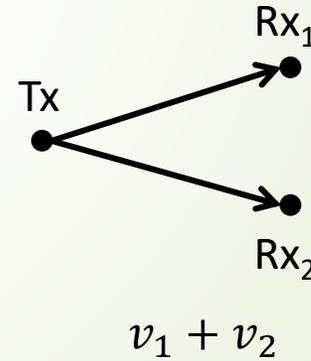
Send user 1's packets



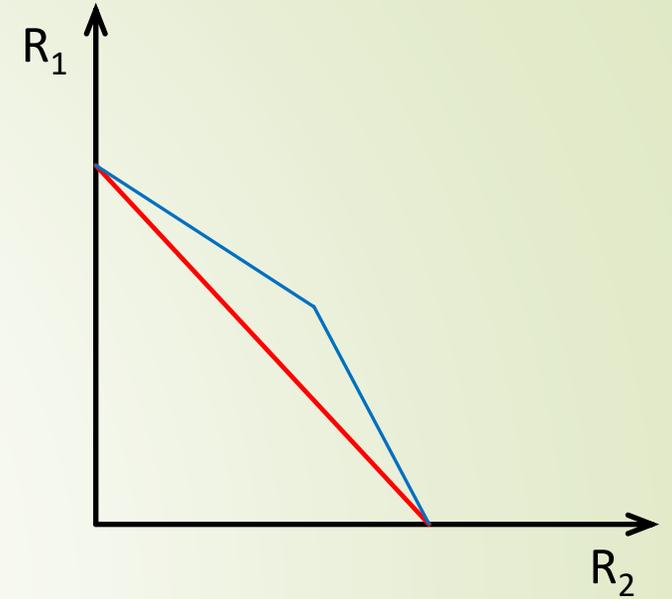
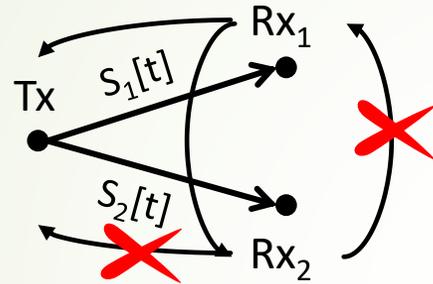
Send user 2's packets



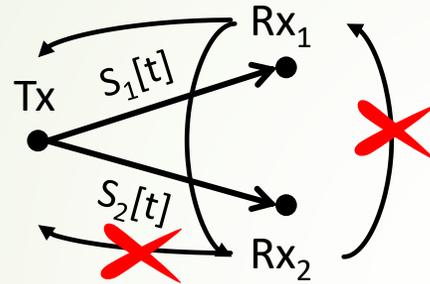
benefit from multicast



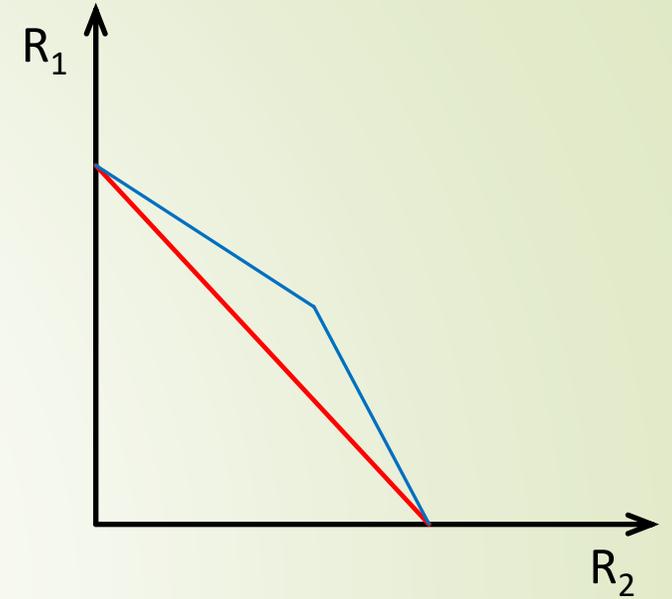
Denial-of-service attack



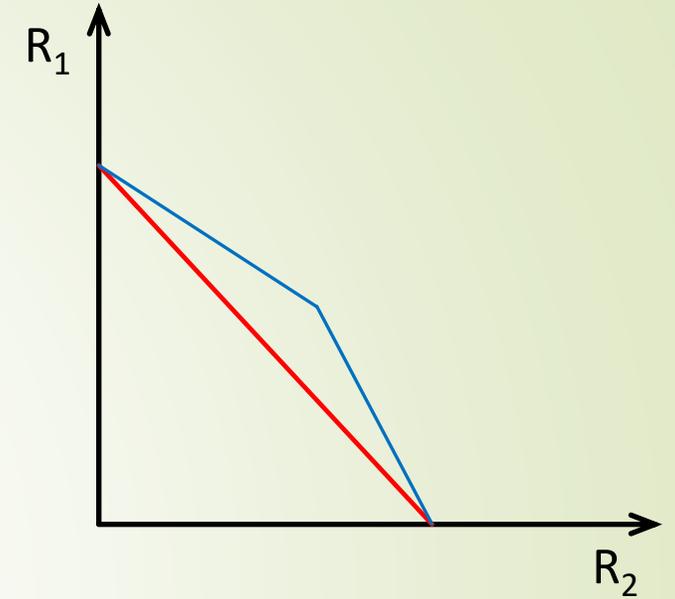
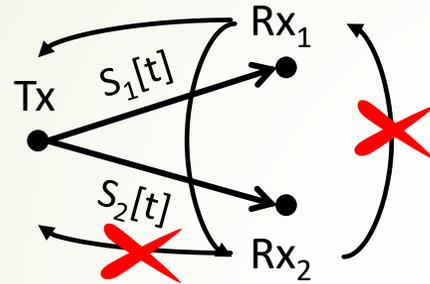
Denial-of-service attack



- Is this single-user knowledge still useful?

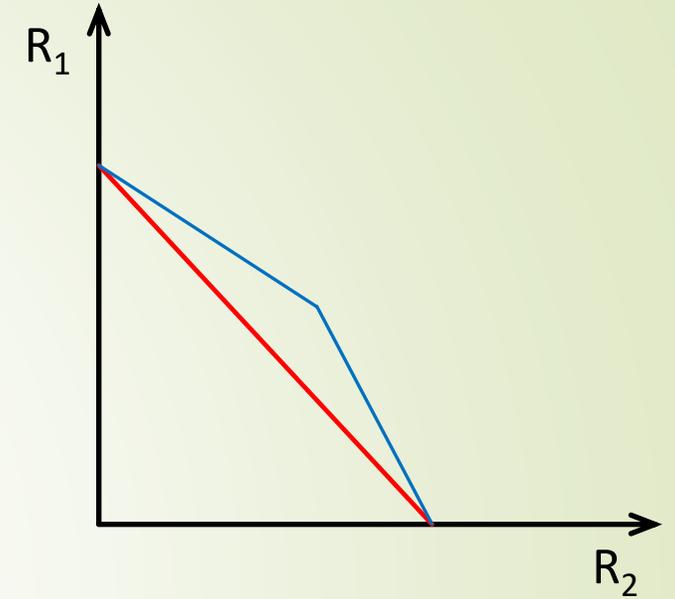
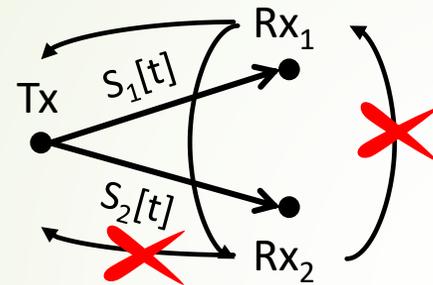


Denial-of-service attack



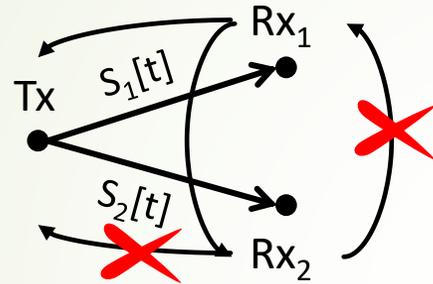
- Is this single-user knowledge still useful?
 - For MISO BC with continuous feedback, the answer is no!

Denial-of-service attack

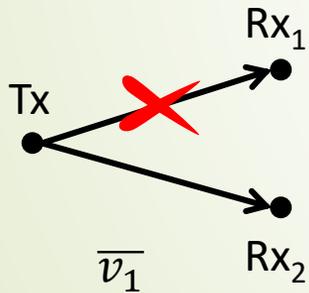


- Is this single-user knowledge still useful?
 - For MISO BC with continuous feedback, the answer is no!
 - We have a much brighter picture in packet networks!

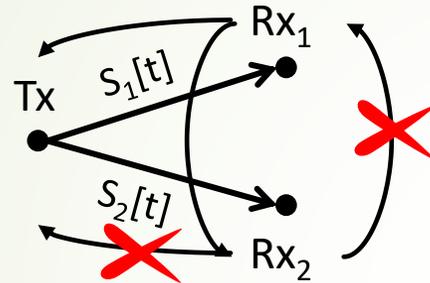
Protocol under strong denial-of-attack



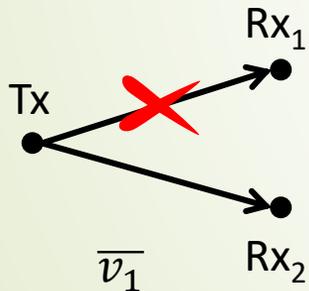
Send user 1's packets



Protocol under strong denial-of-attack

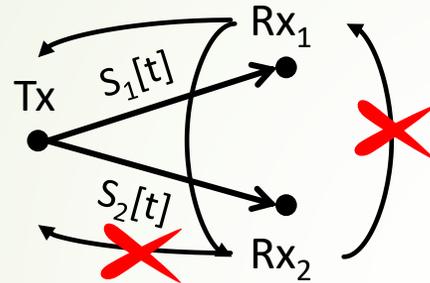


Send user 1's packets

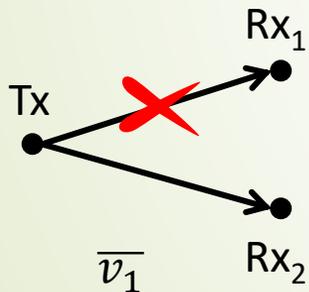


I know what user 1
is missing; and statistically
what user 2 gets,

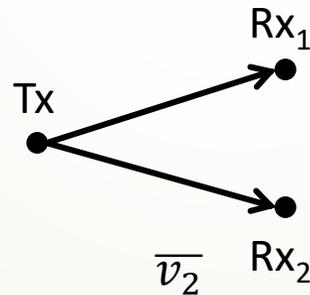
Protocol under strong denial-of-attack



Send user 1's packets

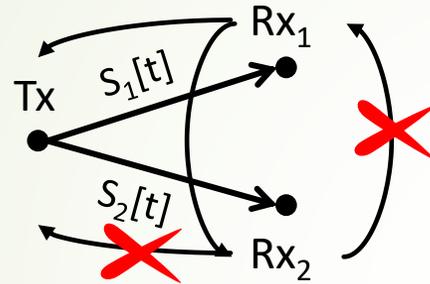


Send user 2's packets

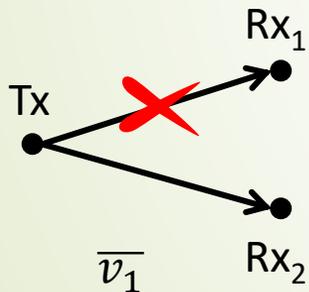


I know what user 1 is missing; and statistically what user 2 gets,

Protocol under strong denial-of-attack

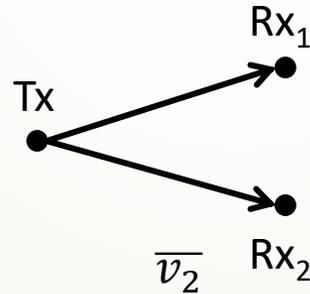


Send user 1's packets



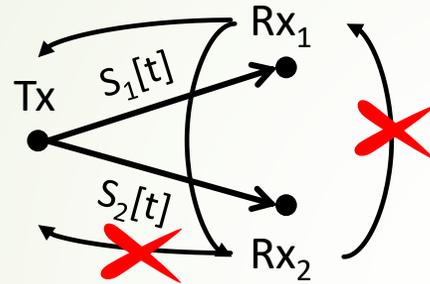
I know what user 1 is missing; and statistically what user 2 gets,

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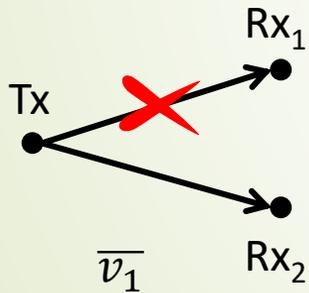


We don't know when user 2 was off! But we know what user 1 receives.

Protocol under strong denial-of-attack

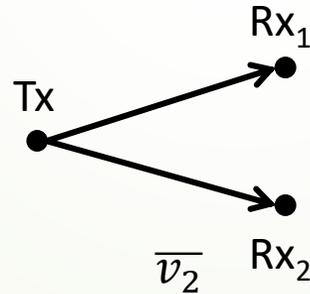


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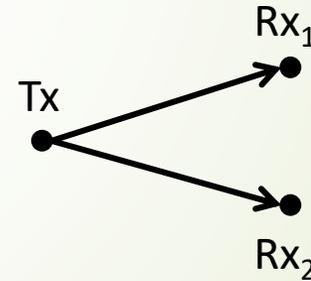
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Send user 2's packets

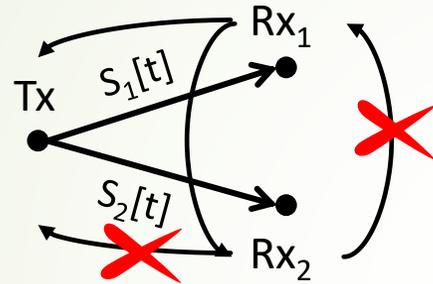


We don't know when user 2 was off! But we know what user 1 receives.

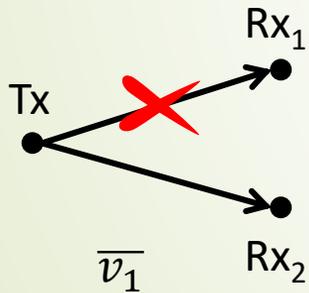
benefit from **feedback**



Protocol under strong denial-of-attack

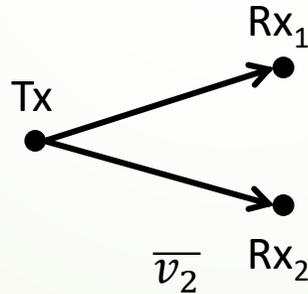


Send user 1's packets



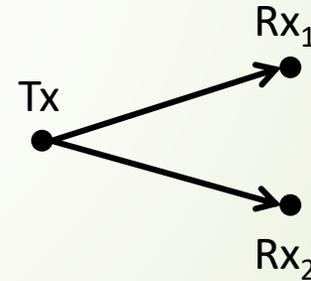
I know what user 1 is missing; and statistically what user 2 gets,

Send user 2's packets



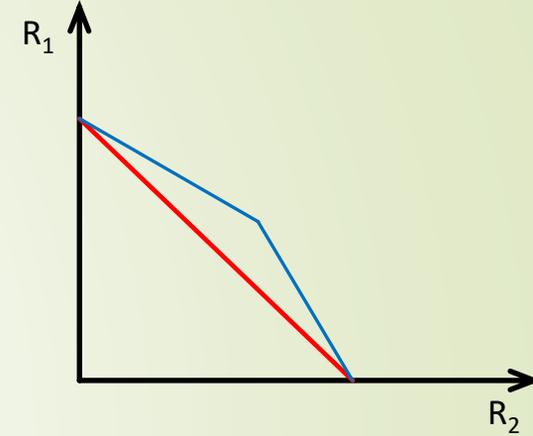
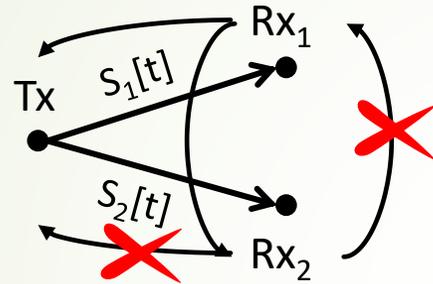
We don't know when user 2 was off! But we know what user 1 receives.

benefit from **feedback**

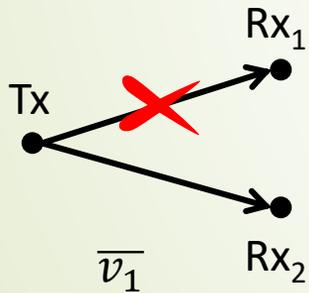


Resend \bar{v}_1 until ACK
+
Linearly coded \bar{v}_2

Protocol under strong denial-of-attack

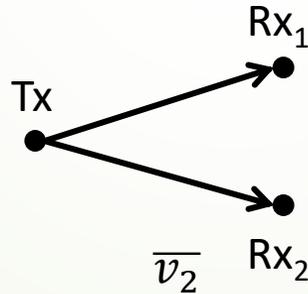


Send user 1's packets



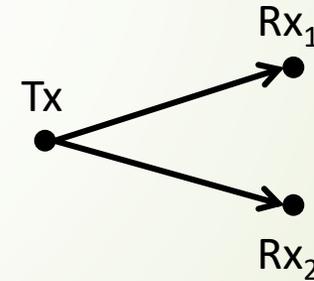
I know what user 1 is missing; and statistically what user 2 gets,

Send user 2's packets



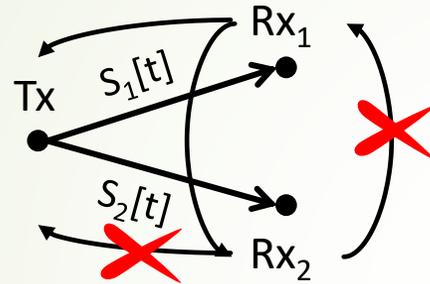
We don't know when user 2 was off! But we know what user 1 receives.

benefit from **feedback**

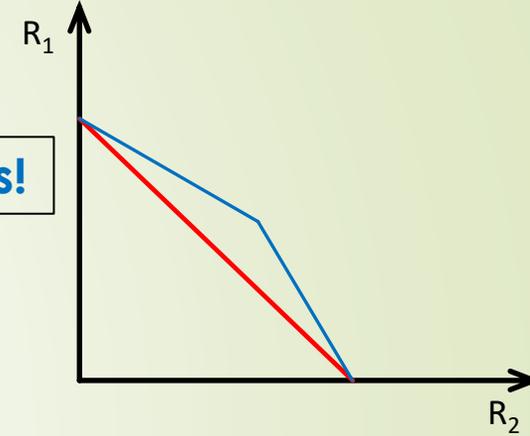


Resend \bar{v}_1 until ACK
+
Linearly coded \bar{v}_2

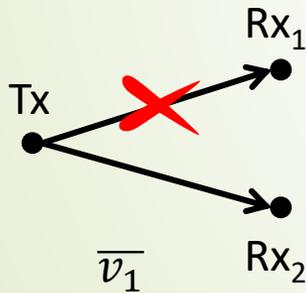
Protocol under strong denial-of-attack



No throughput loss!

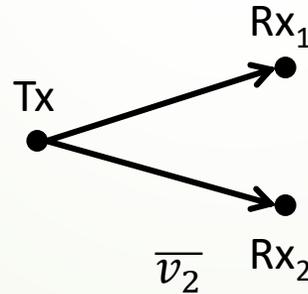


Send user 1's packets



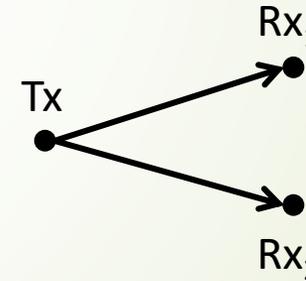
I know what user 1 is missing; and statistically what user 2 gets,

Send user 2's packets



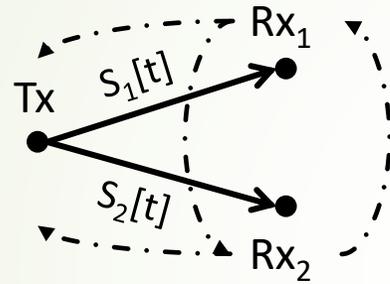
We don't know when user 2 was off! But we know what user 1 receives.

benefit from **feedback**

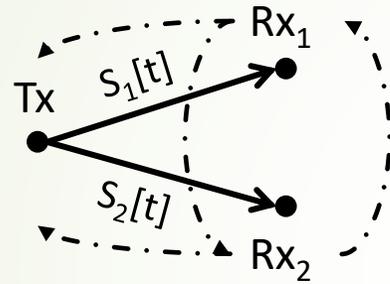


Resend \bar{v}_1 until ACK
+
Linearly coded \bar{v}_2

General denial-of-attack on control channels



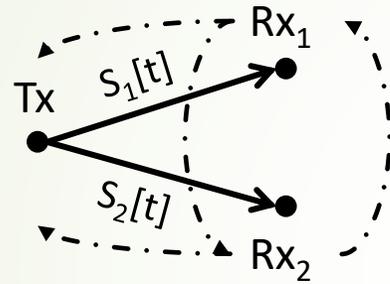
General denial-of-attack on control channels



Each Rx simply broadcasts its control packet.

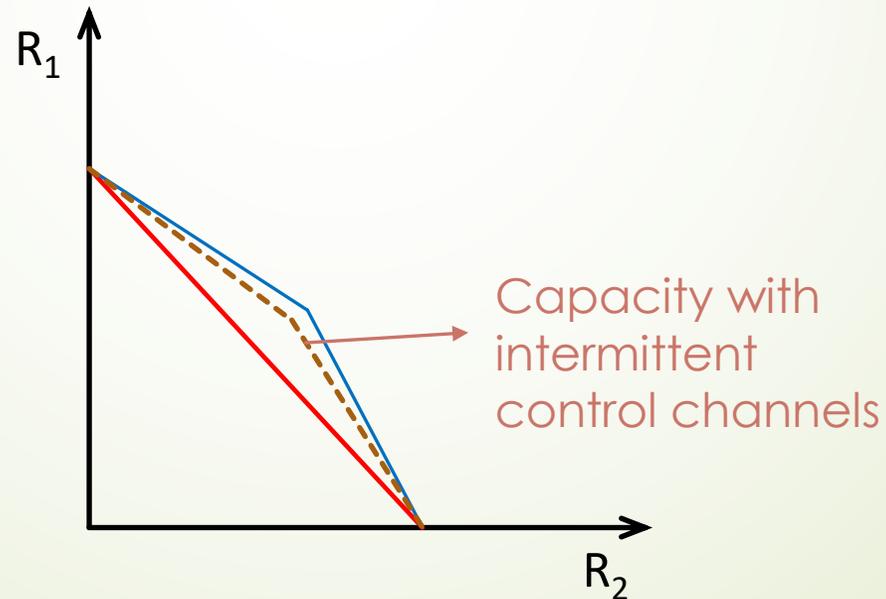
All control channels have some probability of failure.

General denial-of-attack on control channels

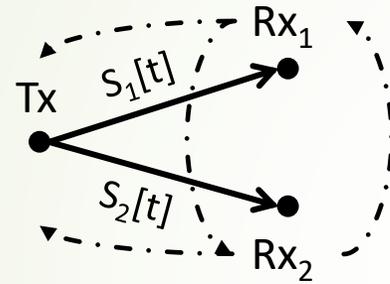


Each Rx simply broadcasts its control packet.

All control channels have some probability of failure.

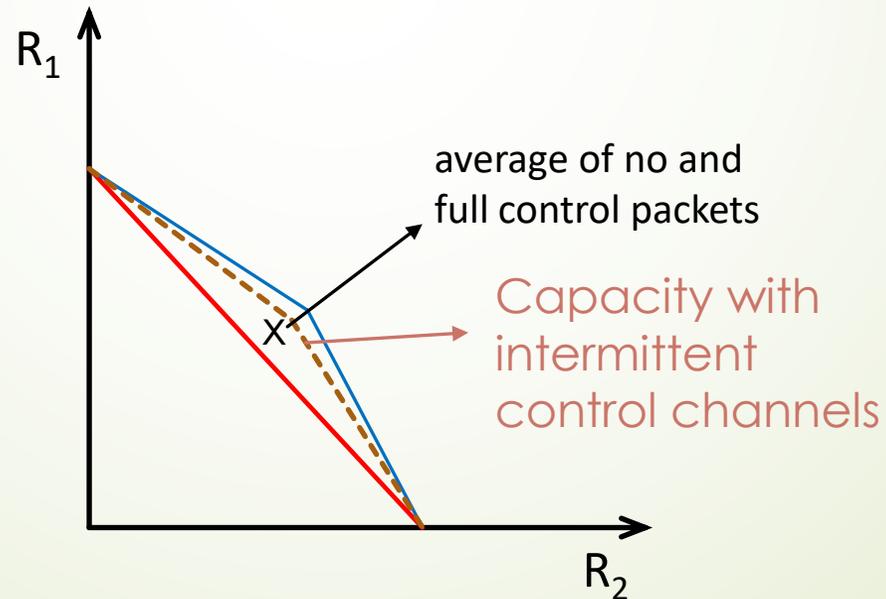


General denial-of-attack on control channels

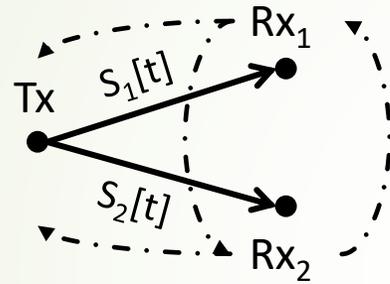


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General denial-of-attack on control channels

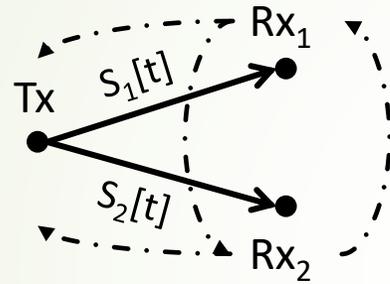


Each Rx simply broadcasts its control packet.

All control channels have some probability of failure.

- ➔ Phase 1: Send bits for user 1.

General denial-of-attack on control channels

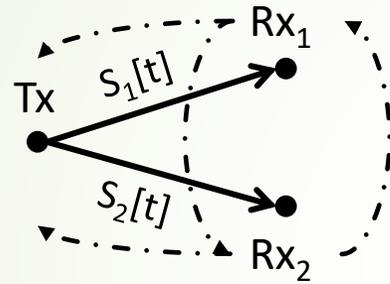


Each Rx simply broadcasts its control packet.

All control channels have some probability of failure.

- Phase 1: Send bits for user 1.
 - when there is FB: v_1 are the bits at Rx₂ needed at Rx₁
 - when no FB: \bar{v}_1 are statistical equations needed at Rx₁

General denial-of-attack on control channels

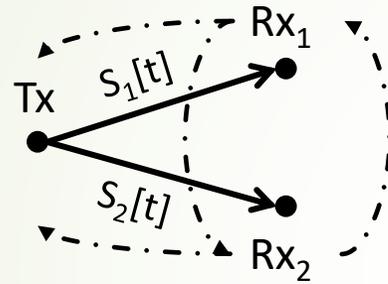


Each Rx simply broadcasts its control packet.

All control channels have some probability of failure.

- Phase 1: Send bits for user 1.
 - when there is FB: v_1 are the bits at Rx_2 needed at Rx_1
 - when no FB: \bar{v}_1 are statistical equations needed at Rx_1
- Phase 2: Send bits for user 2. Create \bar{v}_2 and v_2 .

General denial-of-attack on control channels

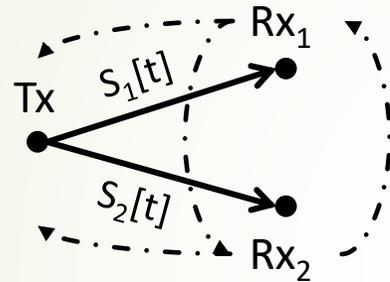


Each Rx simply broadcasts its control packet.

All control channels have some probability of failure.

- Phase 1: Send bits for user 1.
 - when there is FB: v_1 are the bits at Rx_2 needed at Rx_1
 - when no FB: \bar{v}_1 are statistical equations needed at Rx_1
- Phase 2: Send bits for user 2. Create \bar{v}_2 and v_2 .
- Phase 3: send the summation of v_1 & v_2 .

General denial-of-attack on control channels

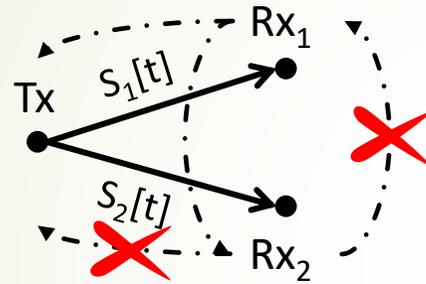


Each Rx simply broadcasts its control packet.

All control channels have some probability of failure.

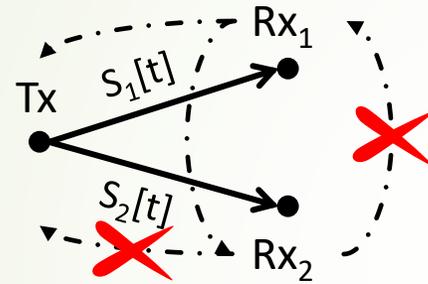
- Phase 1: Send bits for user 1.
 - when there is FB: v_1 are the bits at Rx₂ needed at Rx₁
 - when no FB: \bar{v}_1 are statistical equations needed at Rx₁
- Phase 2: Send bits for user 2. Create \bar{v}_2 and v_2 .
- Phase 3: send the summation of v_1 & v_2 .
- Recursion: Use \bar{v}_1 & \bar{v}_2 as inputs to Phase 1.

Extreme environment

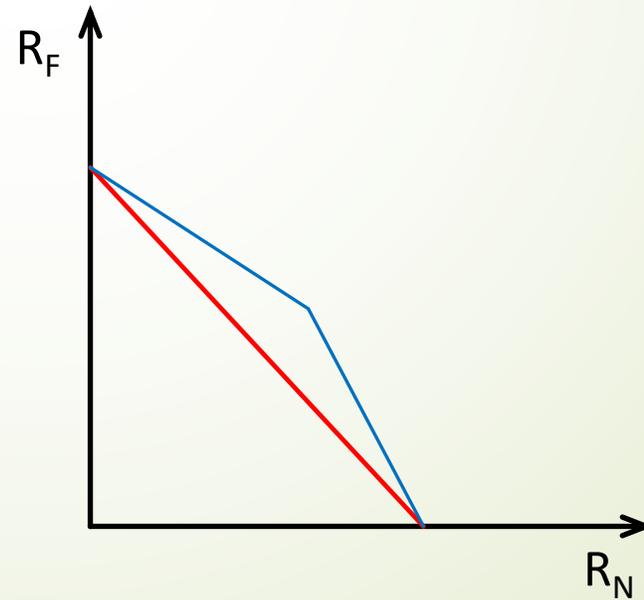


The available control channels have sub-bit capacity!

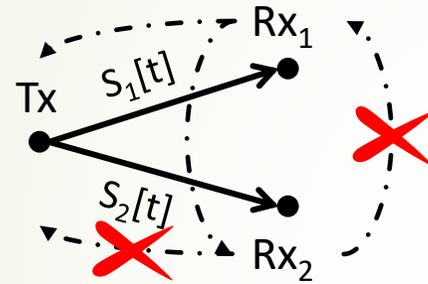
Extreme environment



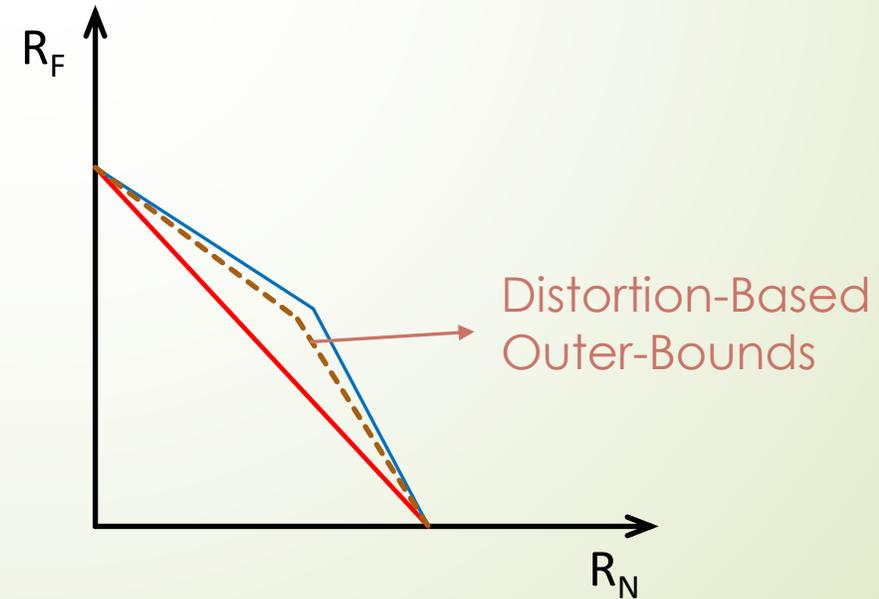
The available control channels have sub-bit capacity!



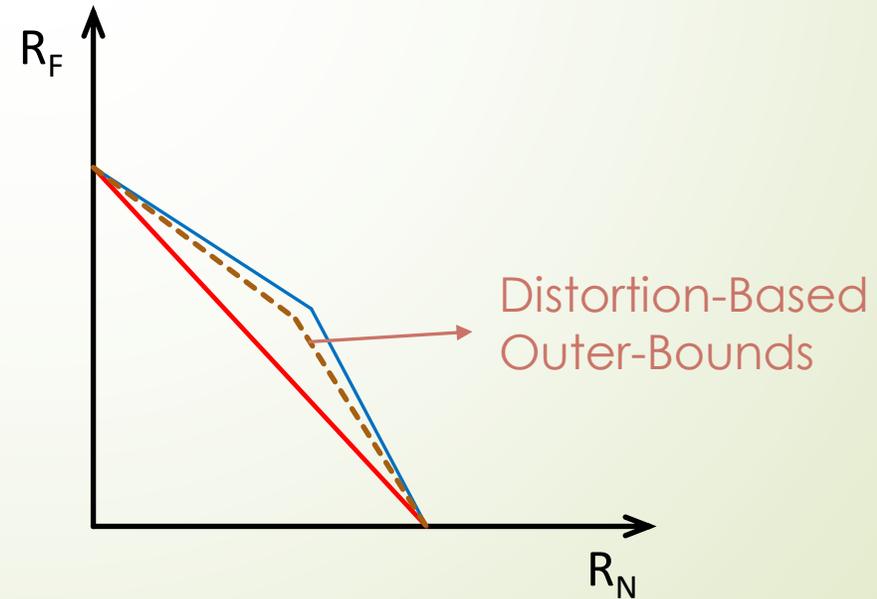
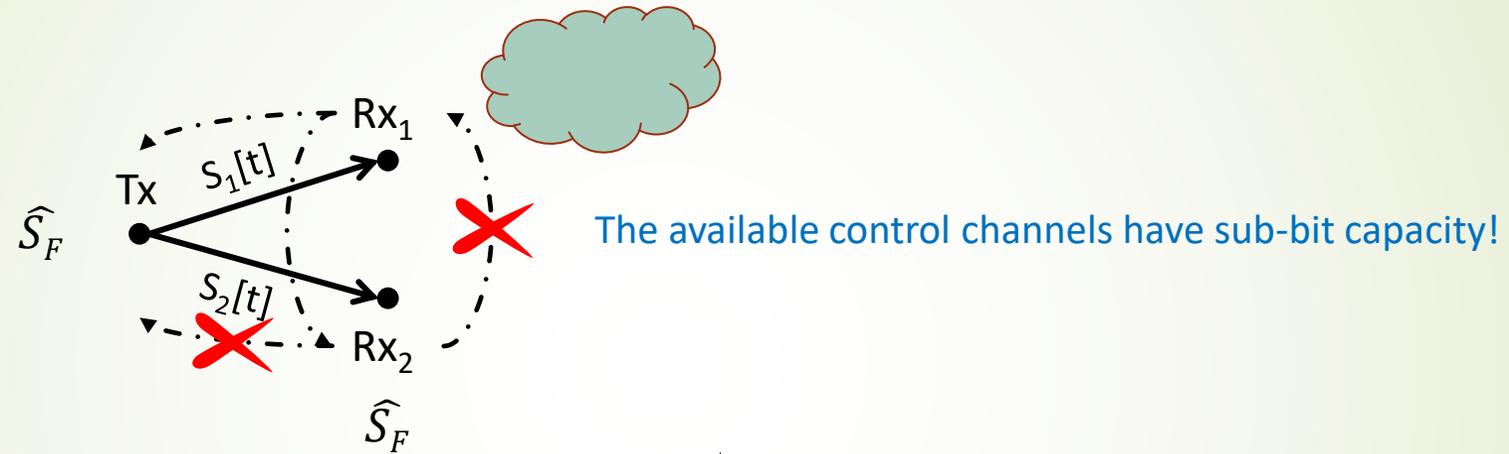
Extreme environment



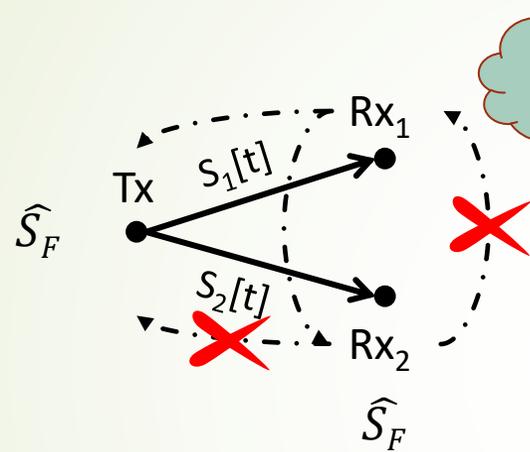
The available control channels have sub-bit capacity!



Extreme environment

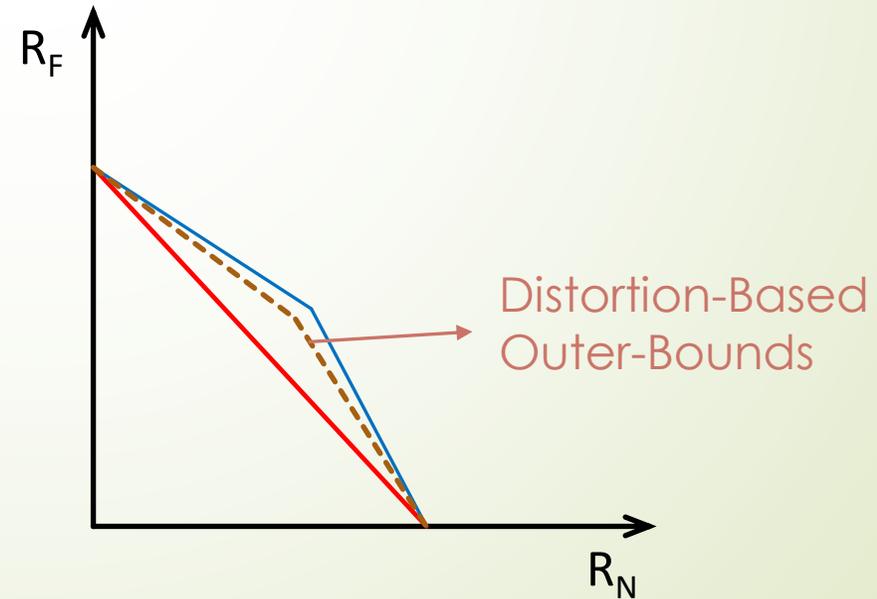


Extreme environment

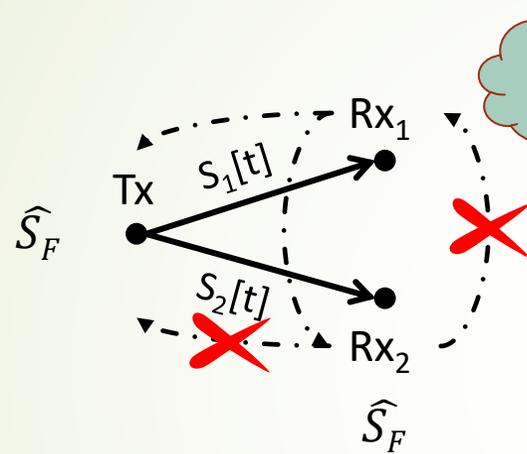


The available control channels have sub-bit capacity!

Rate-distortion theory gives us the minimum attainable distortion.

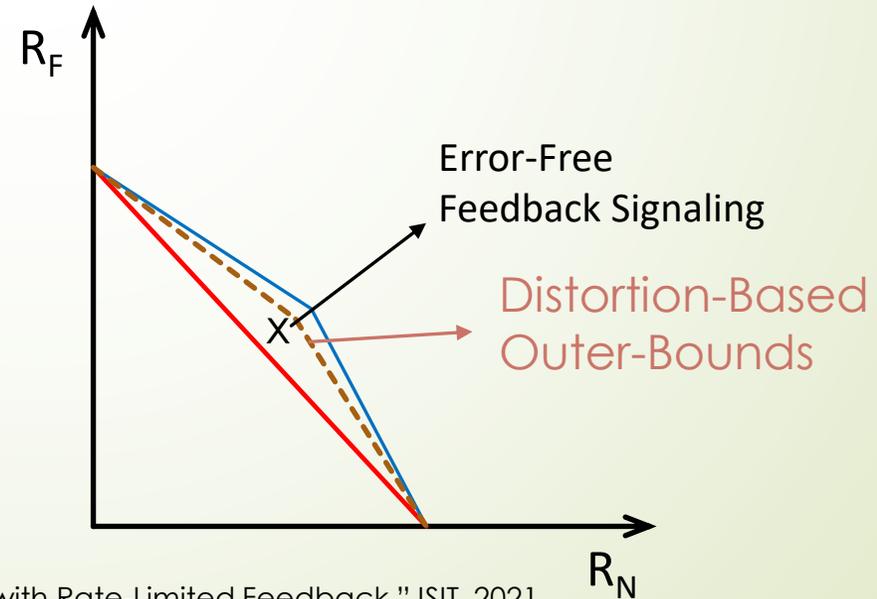


Extreme environment

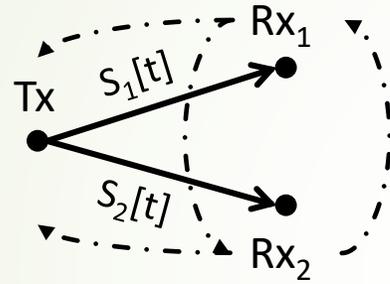


The available control channels have sub-bit capacity!

Rate-distortion theory gives us the minimum attainable distortion.

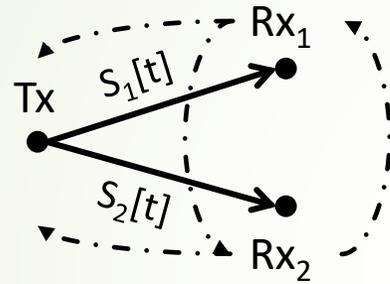


Remaining theoretical questions



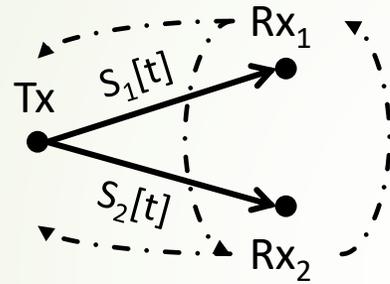
- ▶ How do the results scale?

Remaining theoretical questions



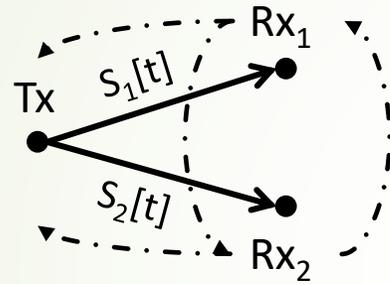
- ▶ How do the results scale?
- ▶ What are the delay implications of the protocols?

Remaining theoretical questions

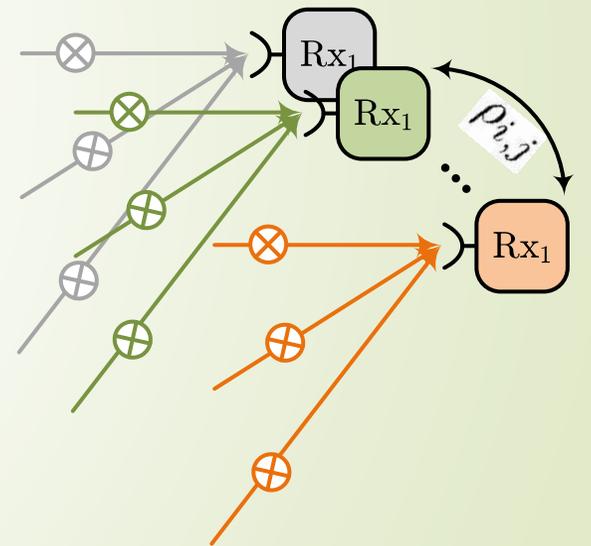


- ▶ How do the results scale?
- ▶ What are the delay implications of the protocols?
- ▶ The extreme sub-bit regime remains open.

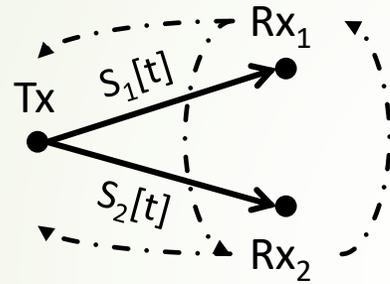
Remaining theoretical questions



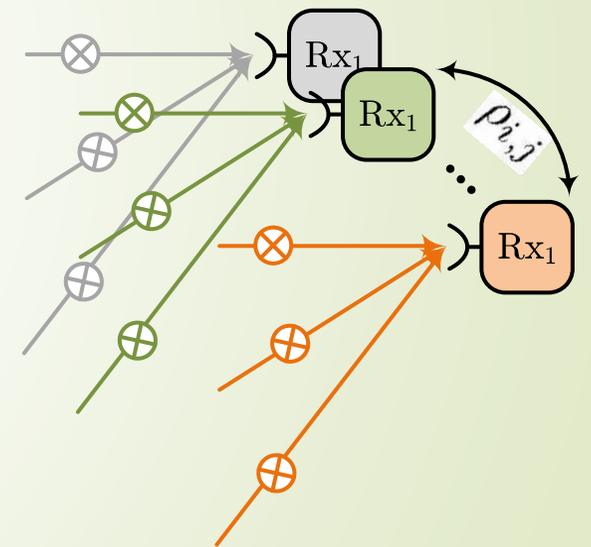
- ▶ How do the results scale?
- ▶ What are the delay implications of the protocols?
- ▶ The extreme sub-bit regime remains open.
- ▶ Spectrum sharing



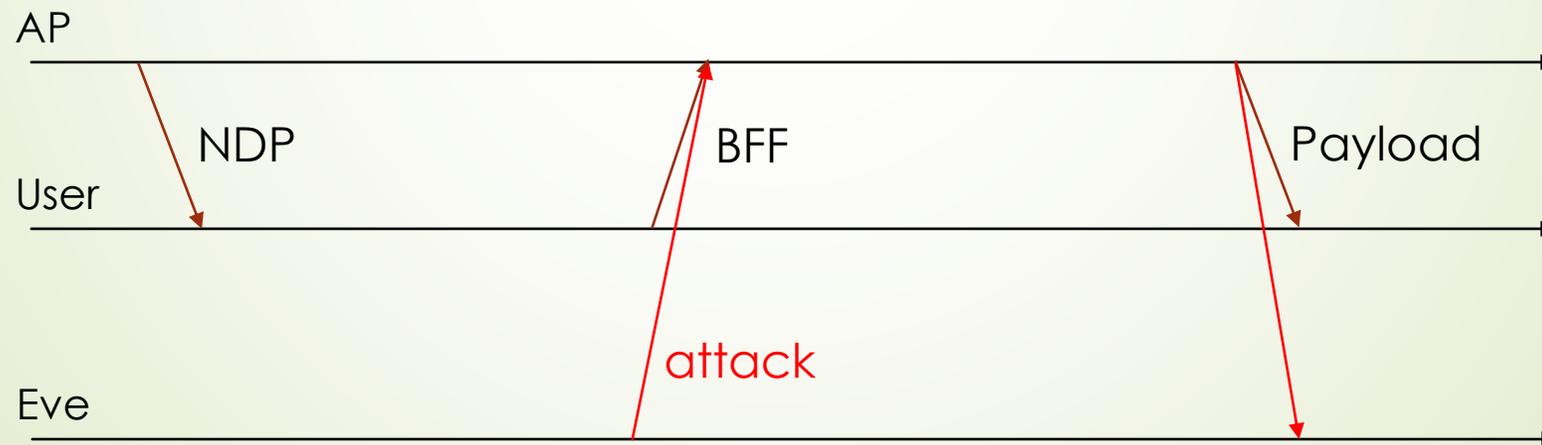
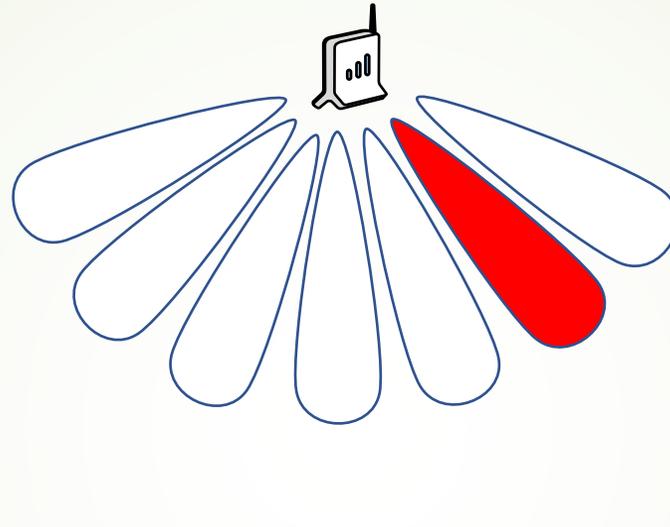
Remaining theoretical questions



- ▶ How do the results scale?
- ▶ What are the delay implications of the protocols?
- ▶ The extreme sub-bit regime remains open.
- ▶ Spectrum sharing
- ▶ Can these ideas be incorporated in existing protocols?



Back to WiFi-6 (-7)

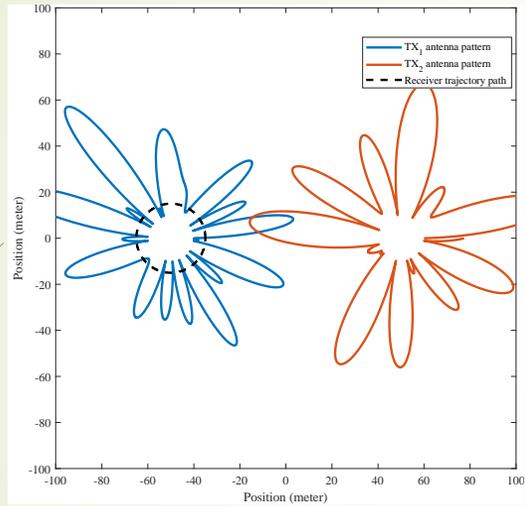




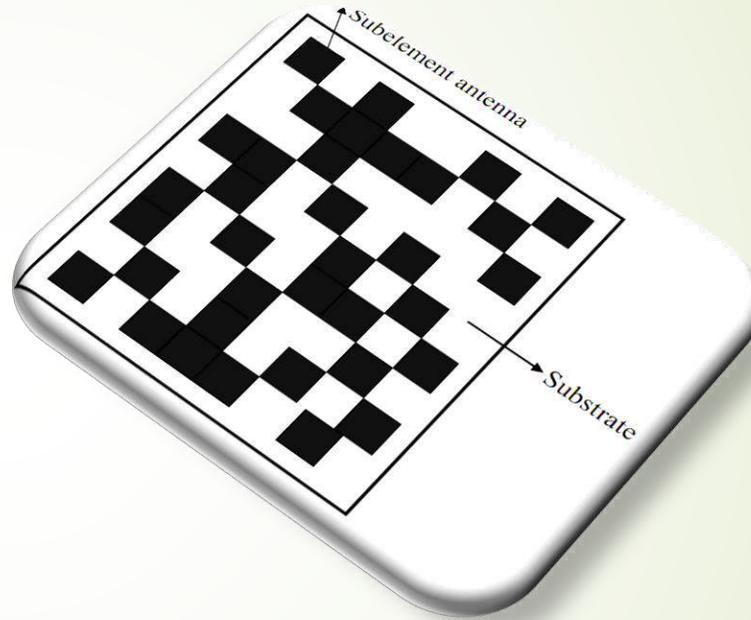
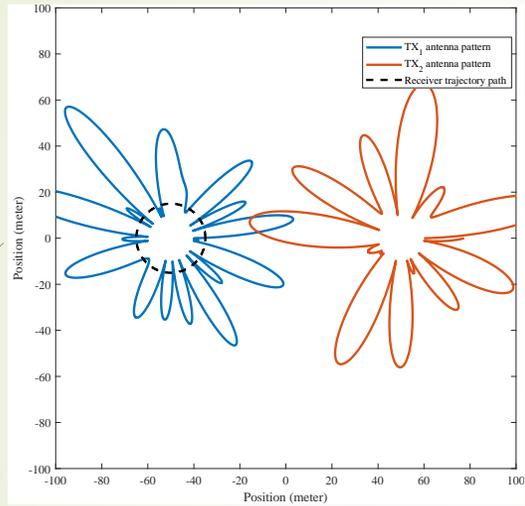
Defending WiFi against Control Ch. Attacks



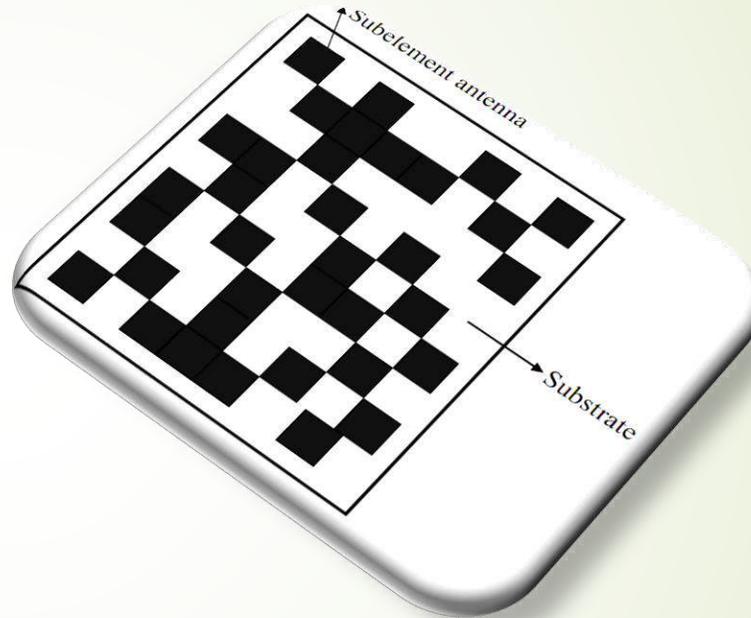
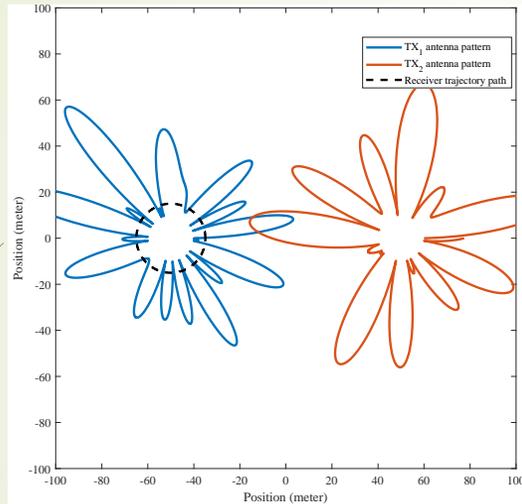
Defending WiFi against Control Ch. Attacks



Defending WiFi against Control Ch. Attacks

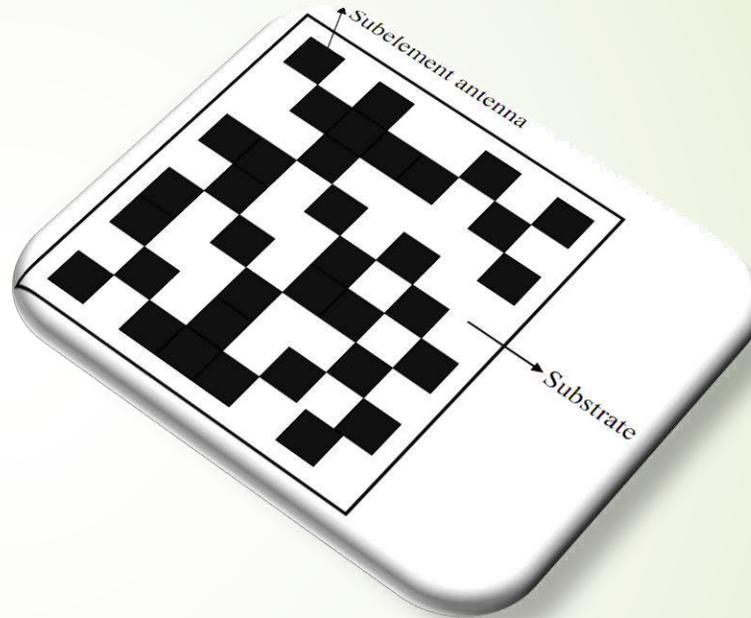
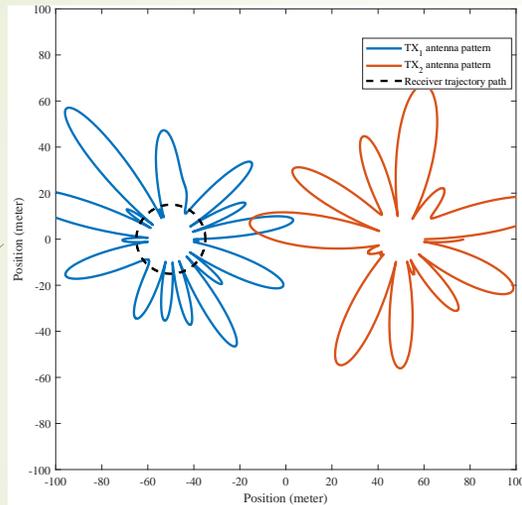


Defending WiFi against Control Ch. Attacks



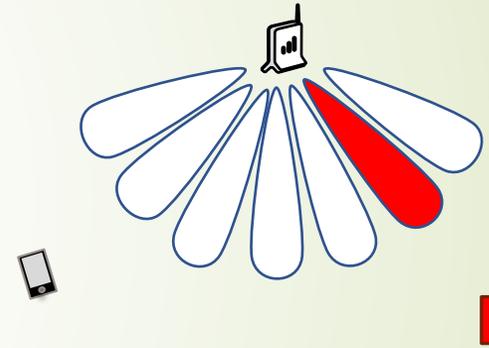
- Desired user will always be at 0° phase, while others see varying phases. (rel. to antenna-selection mod.)

Defending WiFi against Control Ch. Attacks



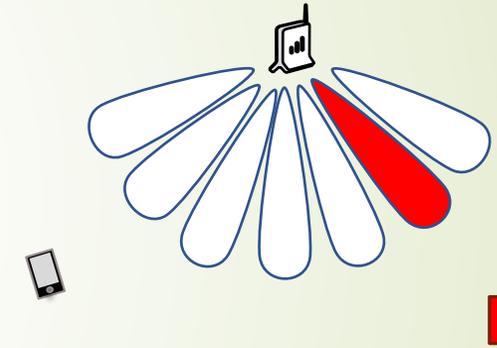
- Desired user will always be at 0° phase, while others see varying phases. (rel. to antenna-selection mod.)
- Embedding information in radiation pattern fluctuations is itself a worthy direction.

Defending WiFi against Control Ch. Attacks



Defending WiFi against Control Ch. Attacks

- Radiation pattern fluctuations.
- WiFi localization (e.g., time of flight).
- Channel signatures.



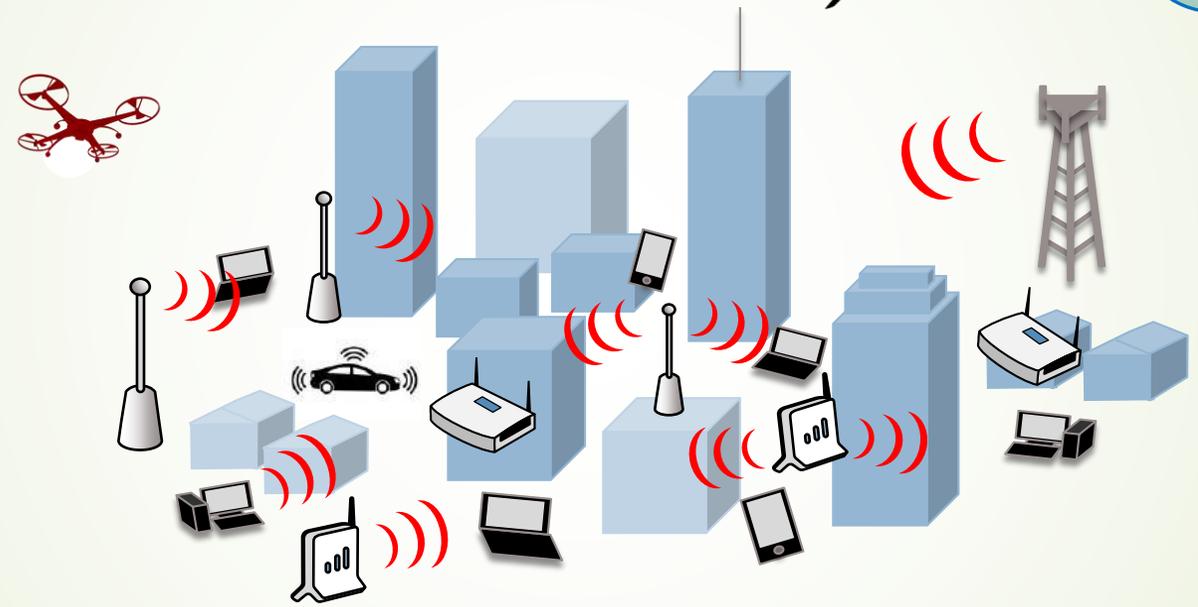
Vahid et al, "A Game-Theoretically Optimal Defense Paradigm against Traffic Analysis Attacks Using Multipath Routing and Deception," SACMAT, 2022.

Vahid et al, "Toward practical defense against traffic analysis attacks on encrypted DNS traffic," Computers & Security, 2022.



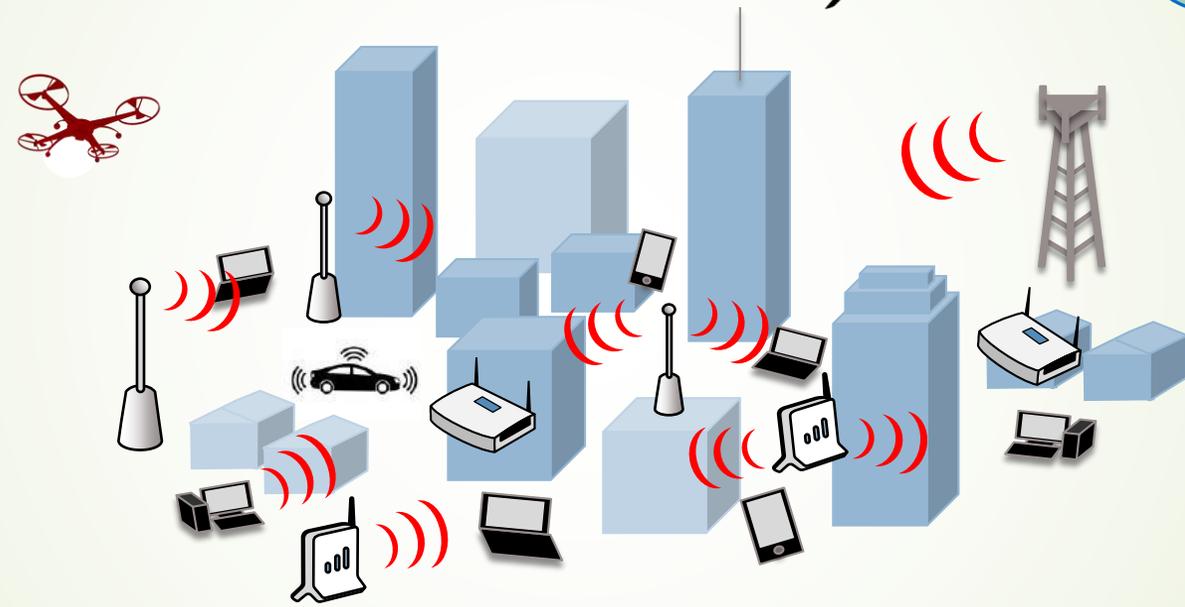


Higher frequency bands?





Higher frequency bands?



Thank you!