

Trust and Resilience in Distributed Consensus Cyberphysical Systems

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October 18, 2021

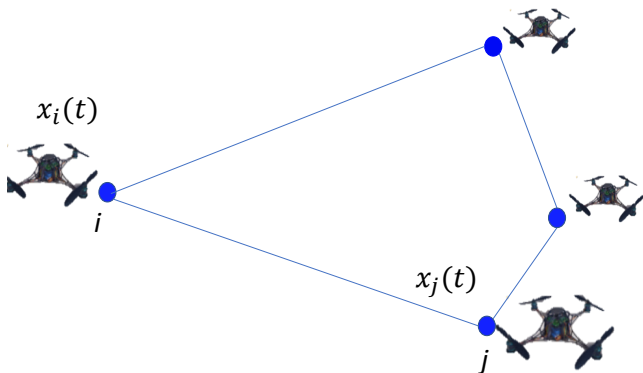
Outline

1. Distributed Consensus Systems.
2. Malicious Agents in Distributed Consensus Systems.
3. Agents' Trust Values in Cyberphysical Systems.
4. Characterizing Trust-Based Resilience in Distributed Consensus Systems.
5. Numerical Results.
6. Conclusions and Future Work.

Distributed Consensus Systems

Leaderless coordination and control for multi-agent systems.

- ▶ Robotic and drone networks (rendezvous problem).
- ▶ Sensor networks (data fusion - temperature measurement).
- ▶ Social networks (reaching a common opinion).



Mathematical of Modeling Distributed Consensus Systems

A **connected** graph $G = (\mathbb{V}, \mathbb{E})$, a **stochastic weight** matrix W and **initial vector** values $x(0)$.

For all $t \geq 0$

$$x_i(t+1) = w_{ii}x_i(t) + \sum_{j \in \mathcal{N}_i} w_{ij}x_j(t),$$

where $\mathcal{N}_i = \{j \in \mathbb{V} \mid \{i, j\} \in \mathbb{E}\}$ and

$$w_{ii} > 0, \quad w_{ij} > 0 \quad \text{for all } j \in \mathcal{N}_i,$$

M. DeGroot 1970's (opinion dynamics), J. Tsitsiklis 1980's (distributed optimization)

It follows that

$$\lim_{t \rightarrow \infty} x_i(t) = \left[\lim_{t \rightarrow \infty} W^t x(0) \right]_i = [\mathbf{1}v'x(0)]_i = \lim_{t \rightarrow \infty} x_j(t), \forall i, j$$

where v' is the Perron-Frobenius left-eigenvector of W .

Malicious Agents in Distributed Consensus Systems

In practice not all agents are legitimate (truthful), some are malicious and strategically input malicious values to either:

- ▶ prevent consensus,
- ▶ deviate the consensus from its true value.

The Classical Bound

The maximal number malicious agents that can be tolerated:

Legitimate agents can reach consensus iff the number of malicious agents is *less than 1/2 of the network connectivity*¹.

Proofs:

- ▶ Lamport, Pease and Shostak 1980, D. Dolev 1981 (Byzantine, fault tolerance, an additional condition),
- ▶ F. Pasqualetti, A. Bicchi and F. Bullo 2012 (control theory).

Both proofs assume that every legitimate agent knows the topology of G , and cannot detect malicious agents that only lie about their initial input values.

¹The connectivity of a graph is the maximum number of disjoint paths between any two vertices of the graph.

Agents' Trust Values in Cyberphysical Systems I

Prior works have used the **data values** to overcome/detect malicious behavior. The **physical** aspects of the problem have not been considered. Namely, the **wireless communication channels**.

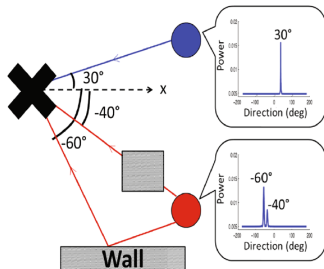
In cyberphysical systems:

- ▶ Malicious agents can lie about their location.
- ▶ A malicious agent can create many fictitious identities (Sybil attack).

Agents' Trust Values in Cyberphysical Systems II

Each transmitted signal leads to a received signal characteristics:

- ▶ Number of paths, delays.
- ▶ Angles of arrival.
- ▶ Power order of the angles of arrival.
- ▶ Power of the received signals.



*Guaranteeing spoof-resistant multi-robot networks, S. Gil *et al* 2017.

Agents' Trust Values in Cyberphysical Systems III

We can generate trust values that captures the event that an agent

- ▶ lies about its location
 - ▶ Location Verification Systems for VANETs in Rician Fading Channels, S. Yan *et al* 2016.
- ▶ uses a Sybil attack and creates multiple fictitious agents
 - ▶ Detecting Colluding Sybil Attackers in Robotic Networks using Backscatters Y. Huang *et al* 2021.
(Limited to single antenna malicious agents.)
 - ▶ Guaranteeing spoof-resilient multi-robot networks, S. Gil *et al* 2017.
(Limited to single antenna malicious agents.)
 - ▶ The Mason Test: A Defense Against Sybil Attacks in Wireless Networks Without Trusted Authorities, Liu *et al* 2015.
(Assumes limited mobility of malicious agents and no beamforming).

We denote by $\alpha_{ij}(t) \in [0, 1]$ the instantaneous single sample trust agent i gives agent j at time a t .

The Trust Based Distributed Consensus Model

Consider the system

$$\begin{bmatrix} X_{\mathcal{L}}(t+1) \\ X_{\mathcal{M}}(t+1) \end{bmatrix} = \begin{bmatrix} W_{\mathcal{L}}(t) & W_{\mathcal{M}}(t) \\ \Theta(t) & \Omega(t) \end{bmatrix} \begin{bmatrix} X_{\mathcal{L}}(t) \\ X_{\mathcal{M}}(t) \end{bmatrix},$$

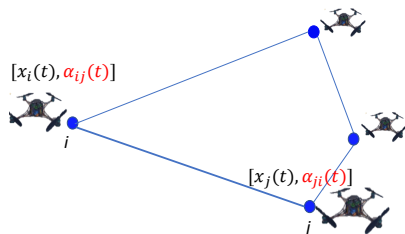
where $|x_i(t)| \leq \eta$ for every $i, j \in \mathcal{L} \cup \mathcal{M}$ and $t \geq 0$.

For every $i \in \mathcal{L}$:

$$x_i(t+1) = \underbrace{\left[1 - \sum_{j \in \mathcal{N}_i} W(i, j, t, \beta_{ij}(t))\right]}_{w_{ii}(t)} x_i(t) + \sum_{j \in \mathcal{N}_i} \underbrace{W(i, j, t, \beta_{ij}(t))}_{w_{ij}(t)} x_j(t)$$

where

► $\beta_{ij}(t) = f(\alpha_{ij}(0), \dots, \alpha_{ij}(t))$



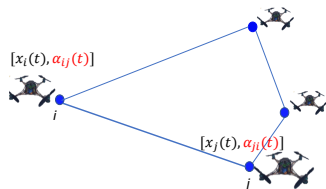
The Trust Based Distributed Consensus Model

For every $i \in \mathcal{L}$:

$$x_i(t+1) = \underbrace{\left[1 - \sum_{j \in \mathcal{N}_i} W(i, j, t, \beta_{ij}(t))\right]}_{w_{ii}(t)} x_i(t) + \sum_{j \in \mathcal{N}_i} \underbrace{W(i, j, t, \beta_{ij}(t))}_{w_{ij}(t)} x_j(t)$$

where

- ▶ $\beta_{ij}(t) = f(\alpha_{ij}(0), \dots, \alpha_{ij}(t))$
- ▶ $w_{ii}(t) > 0$, $w_{ij}(t) \geq 0$, $j \in \mathcal{N}_i$, $\sum_{j \in \mathcal{N}_i} w_{ij} = 1$
- ▶ $w_{ij}(t) > 0$, $j \in \mathcal{N}_i \cap \mathcal{M}$ finitely many times a.s.
- ▶ $w_{ij}(t) = 0$, $j \in \mathcal{N}_i \cap \mathcal{L}$ finitely many times a.s.



Research Objectives

Objective I - Finite correct classification time

Establish characteristics of $\alpha_{ij}(t)$, and functions $\beta_{ij}(t)$ that lead to a **finite detection time** for the correct classification of legitimate and malicious agents **almost surely**.

Objective II - Convergence of the consensus protocol

Choose weights $W(i, j, t, \beta_{ij}(t))$ that allow **convergence** in spite of the presence of adversarial attacks.

Objective III - Bounded deviation for average consensus

We bound the **deviation** from the **true** consensus value, $\Delta(\delta)$ that can be achieved with a probability at least $1 - \delta$.

Cumulative Trust Values

We assume that:

- ▶ $\alpha_{ij}(t)$ are statistically independent.
- ▶ There exist scalars $c < 0$ and $d > 0$ such that²

$$c = c_{ij} = E(\alpha_{ij}(t)) - 1/2 \quad \text{for all } i \in \mathcal{L}, j \in \mathcal{N}_i \cap \mathcal{M},$$

$$d = d_{ij} = E(\alpha_{ij}(t)) - 1/2 \quad \text{for all } i \in \mathcal{L}, j \in \mathcal{N}_i \cap \mathcal{L}.$$

To capture the **history** of observations $\alpha_{ij}(t)$, we define:

$$\beta_{ij}(t) = \sum_{k=0}^t (\alpha_{ij}(k) - 1/2) \quad \text{for } t \geq 0, i \in \mathcal{L}, j \in \mathcal{N}_i.$$

Agent i classifies agent j as legitimate if $\beta_{ij}(t) \geq 0$ and malicious otherwise.

²For the sake of simplicity of presentation.

Finite Correct Classification Time I

Lemma

For every $t \geq 0$ and $i \in \mathcal{L}$

$$\Pr(\beta_{ij}(t) < 0) \leq \exp(-2(t+1)d^2), j \in \mathcal{N}_i \cap \mathcal{L},$$

$$\Pr(\beta_{ij}(t) \geq 0) \leq \exp(-2(t+1)c^2), j \in \mathcal{N}_i \cap \mathcal{M}.$$

This is an immediate result of the Chernoff-Hoeffding Inequality.

Proposition

There exists a (random) finite time instant $T_f > 0$ such that every legitimate agent i correctly classifies its neighbors for all $t \geq T_f$ almost surely.

This proposition follows by the Borel-Cantelli Lemma

The Modified Trust Based Weights

Define the time dependent **trusted neighborhood** for agent i :

$$\mathcal{N}_i(t) = \{j \in \mathcal{N}_i : \beta_{ij}(t) \geq 0\},$$

We choose for all $i \in \mathcal{L}$,

$$w_{ij}(t) = \begin{cases} \mathbb{1}_{\{t \geq T_0 - 1\}} \cdot \min \left\{ \frac{1}{\kappa}, \frac{1}{|\mathcal{N}_i(t)| + 1} \right\} & \text{if } j \in \mathcal{N}_i(t), \\ 0 & \text{if } j \notin \mathcal{N}_i(t) \cup \{i\}, \\ 1 - \sum_{m \in \mathcal{N}_i} w_{im}(t) & \text{if } j = i. \end{cases}$$

where $\kappa > 0$ is a limiting effect constant.

Up to time T_0 agents measure the trust values of their neighbors but don't update their data values.

The Data Values of the Legitimate Agents

Recall that:

$$\begin{bmatrix} X_{\mathcal{L}}(t+1) \\ X_{\mathcal{M}}(t+1) \end{bmatrix} = \begin{bmatrix} W_{\mathcal{L}}(t) & W_{\mathcal{M}}(t) \\ \Theta(t) & \Omega(t) \end{bmatrix} \begin{bmatrix} X_{\mathcal{L}}(t) \\ X_{\mathcal{M}}(t) \end{bmatrix}.$$

Thus,

$$x_{\mathcal{L}}(t) = \tilde{x}_{\mathcal{L}}(t) + \phi_{\mathcal{M}}(t),$$

where³

$$\tilde{x}_{\mathcal{L}}(t) = \left(\prod_{k=T_0-1}^{t-1} W_{\mathcal{L}}(k) \right) x_{\mathcal{L}}(0),$$

and

$$\phi_{\mathcal{M}}(t) = \sum_{k=T_0-1}^{t-1} \left(\prod_{l=k+1}^{t-1} W_{\mathcal{L}}(l) \right) W_{\mathcal{M}}(k) x_{\mathcal{M}}(k).$$

³Note that $W_{\mathcal{L}}(k)$ can be substochastic.

Convergence of the Consensus Protocol I

Define a matrix $\overline{W}_{\mathcal{L}}$ such that for every $i, j \in \mathcal{L}$,

$$[\overline{W}_{\mathcal{L}}]_{ij} = \begin{cases} \min \left\{ \frac{1}{\kappa}, \frac{1}{|\mathcal{N}_i|+1} \right\} & \text{if } j \in \mathcal{N}_i \cap \mathcal{L}, \\ 1 - \min \left\{ \frac{|\mathcal{N}_i \cap \mathcal{L}|}{\kappa}, \frac{|\mathcal{N}_i \cap \mathcal{L}|}{|\mathcal{N}_i|+1} \right\} & \text{if } j = i, \\ 0 & \text{otherwise.} \end{cases}$$

Then, almost surely there exists a (random) finite time T_f such that

$$\prod_{k=T_0-1}^{\infty} W_{\mathcal{L}}(k) = \underbrace{\lim_{k \rightarrow \infty} \overline{W}_{\mathcal{L}}^{k - \max\{T_f, T_0\}}}_{\mathbf{1}v'} \prod_{k=T_0-1}^{\max\{T_f, T_0\}-1} W_{\mathcal{L}}(k),$$

and $W_{\mathcal{M}}(t) = \mathbf{0}$ for every $t > T_f$.

Convergence of the Consensus Protocol II

Proposition

Almost surely, there exists a random variable $z(T_0)$ such that

$$\lim_{t \rightarrow \infty} x_{\mathcal{L}}(t) = z(T_0)\mathbf{1},$$

where $z(T_0)$ is in the convex hull of the initial values $x_i(0)$, $i \in \mathcal{L} \cup \mathcal{M}$, and its distribution depends on the starting time T_0 of the data passing phase.

The Deviation from Nominal Consensus Value

Theorem

Given an error level $\delta > 0$, we have the following result

$$\Pr\left(\max_{i \in \mathcal{L}} \limsup_{t \rightarrow \infty} |[x_{\mathcal{L}}(t) - \mathbf{1}v'x_{\mathcal{L}}(0)]_i| \leq \Delta_{\max}(T_0, \delta)\right) \geq 1 - \delta,$$

where $\Delta_{\max}(T_0, \delta) = 2[\tilde{g}_{\mathcal{L}}(T_0, \delta) + \tilde{g}_{\mathcal{M}}(T_0, \delta)]$,

$$\tilde{g}_{\mathcal{L}}(\delta) = \frac{\eta|\mathcal{L}|^2}{\delta} \cdot \frac{\exp(-2T_0d^2)}{1 - \exp(-2d^2)} + \frac{\eta|\mathcal{L}||\mathcal{M}|}{\delta} \cdot \frac{\exp(-2T_0c^2)}{1 - \exp(-2c^2)},$$

and

$$\tilde{g}_{\mathcal{M}}(T_0, \delta) = \frac{\eta|\mathcal{L}||\mathcal{M}|}{\delta \cdot \kappa} \cdot \frac{\exp(-2T_0c^2)}{1 - \exp(-2c^2)}.$$

$$x_{\mathcal{L}}(t) = \tilde{x}_{\mathcal{L}}(t) + \phi_{\mathcal{M}}(t) \Rightarrow$$

$$|x_{\mathcal{L}}(t) - \mathbf{1}v'x_{\mathcal{L}}(0)|_i \leq |\tilde{x}_{\mathcal{L}}(t) - \mathbf{1}v'x_{\mathcal{L}}(0)|_i + |\phi_{\mathcal{M}}(t)|_i.$$

A Few Words regarding the Expected Convergence Time I

Proposition

Assume that $j \in \mathcal{N}_i \Leftrightarrow i \in \mathcal{N}_j$ (symmetric connectivity of legitimate agents). Then, for every $T_0 \geq 0$ and $t \geq T_0$, we have

$$\begin{aligned} & E \left(\|x_{\mathcal{L}}(t) - \mathbf{1}v'x_{\mathcal{L}}(0)\|_v \right) \\ & \leq 2 \left(\frac{t - T_0}{2} + 1 \right) \rho_2^{\frac{t-T_0}{2}} \eta + \left(\frac{|\mathcal{L}|^2 \exp(-(t + T_0 + 2)d^2)}{1 - \exp(-2d^2)} \right. \\ & \quad \left. + \frac{|\mathcal{L}||\mathcal{M}| \exp(-(t + T_0 + 2)c^2)}{1 - \exp(-2c^2)} \right) 2\eta \\ & = O \left(|\mathcal{L}| \cdot \max \{ |\mathcal{L}|, |\mathcal{M}| \} \cdot t e^{-\gamma t} \right), \end{aligned}$$

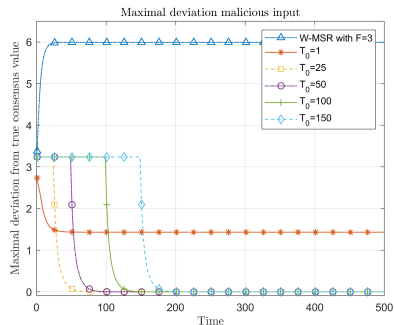
where $\rho_2 < 1$ is the second largest eigenvalue modulus of $\overline{W}_{\mathcal{L}}$ and $v > \mathbf{0}$ be the stochastic Perron vector satisfying $v' \overline{W}_{\mathcal{L}} = v'$.

Numerical Results

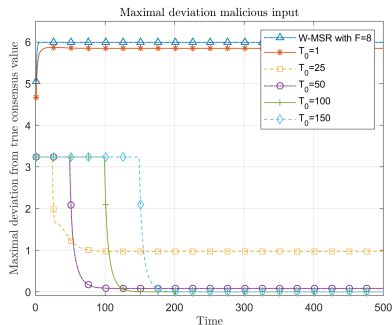
- ▶ $|\mathcal{L}| = 15$ legitimate agents
- ▶ $|\mathcal{M}| = 5, 15, 30$
- ▶ $\eta = 5, \kappa = 10$;
- ▶ $E(\alpha_{ij}) = 0.55$ for $i \in \mathcal{L}, j \in \mathcal{N}_i \cap \mathcal{L}$,
- ▶ $E(\alpha_{ij}) = 0.45$ for $i \in \mathcal{L}, j \in \mathcal{N}_i \cap \mathcal{M}$,
- ▶ $\alpha_{ij} \sim U \left[E(\alpha_{ij}) - \frac{\ell}{2}, E(\alpha_{ij}) + \frac{\ell}{2} \right]$
- ▶ $\ell = 0.2, 0.4, 0.6$

Classical bound must fulfill $|\mathcal{M}| < \frac{3+|\mathcal{M}|}{2} \Rightarrow |\mathcal{M}| < 3$.

Numerical Results - Maximum Deviation Input I

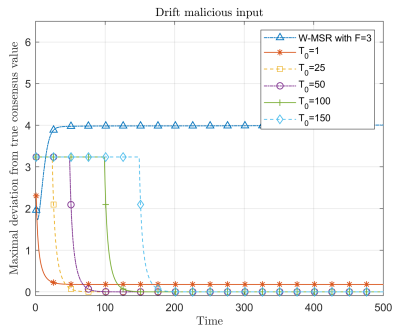


(a) $|\mathcal{M}|=5$, $\ell=0.2$

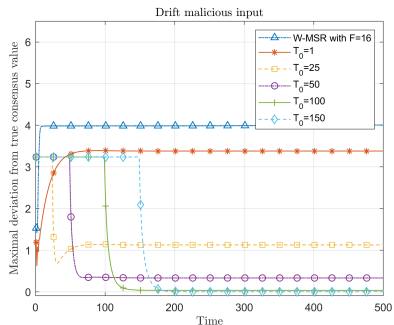


(b) $|\mathcal{M}|=15$, $\ell=0.4$

Numerical Results - Drift Input I



(a) $|\mathcal{M}|=5$, $\ell=0.2$



(b) $|\mathcal{M}|=30$, $\ell=0.6$

Conclusions and Future Work

- ▶ Physical based trust values to brake the current known bound
- ▶ Modified weight matrix - based on trust values
- ▶ Finite detection time a.s., convergence, deviation from true consensus value
- ▶ Future work

Thank You!

Questions? Collaboration ideas?
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